



## Recent algorithmic advances for combinatorial optimization in graphical models

David Allouche, Simon de Givry, Georgios Katsirelos, Thomas Schiex,  
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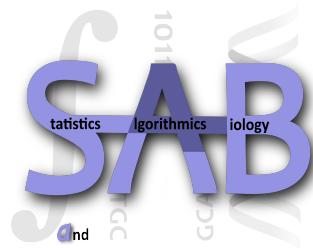
# ISMP'18



## Recent algorithmic advances for combinatorial optimization in graphical models

Simon de Givry, Thomas Schiex, David Allouche, George Katsirelos,  
Matthias Zytnicki, MIAT – INRA, Toulouse, France

Abdelkader Ouali, Samir Loudni, GREYC, University of Caen, France



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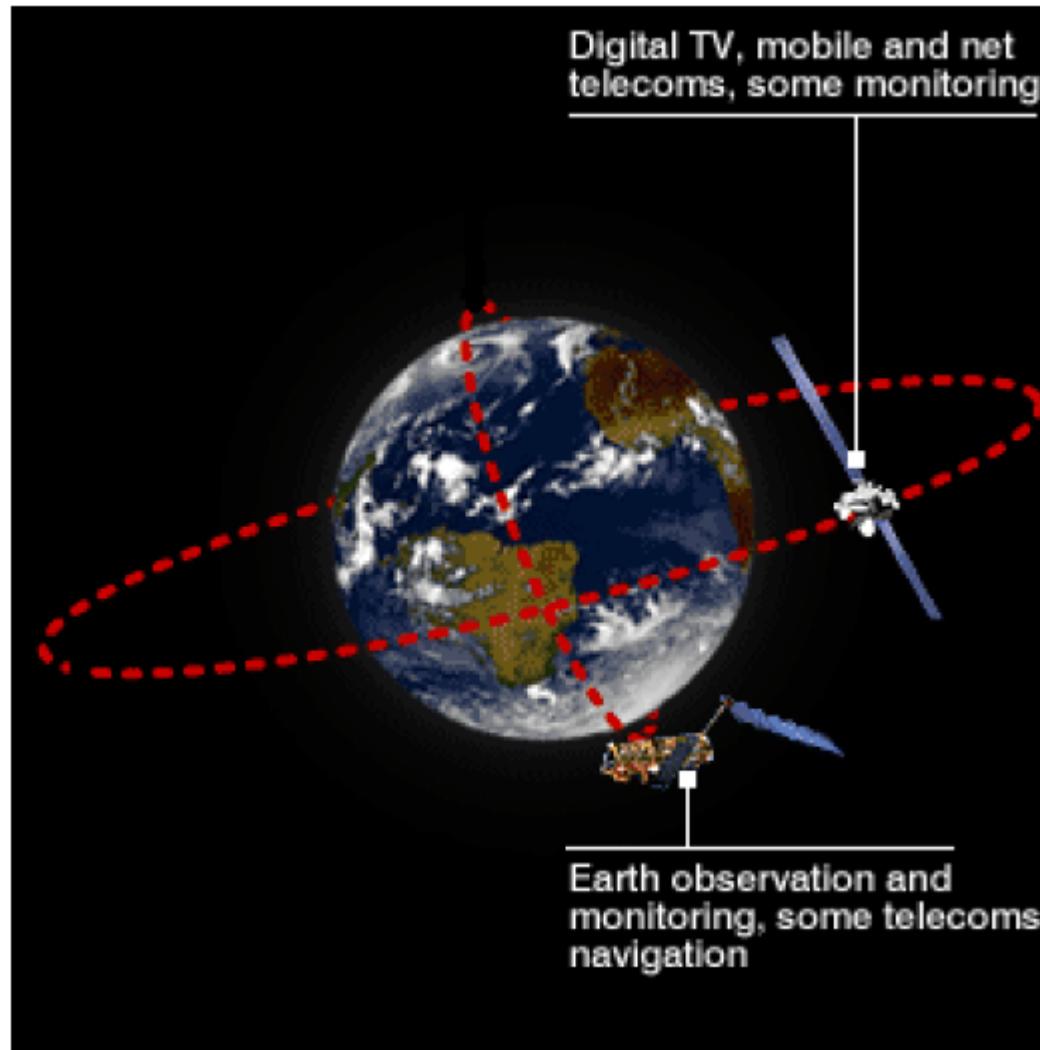
# Plan

- Graphical models
  - Examples and definitions
- Local reasoning techniques
  - Bounding, clique cut, pruning
- Complete search methods
  - Hybrid search, iterative search, large neighborhood search
- Experimental results
  - Open-source C++ exact solver **toulbar2 v1.0.0**

<https://github.com/toulbar2/toulbar2>



# Earth Observation Satellite Management (SPOT5)



(Bensana et al, Constraints 1999 ; IJCAI09)

$n \leq 364$ ,  $d=4$ ,  $e(2-3) \leq 10,108$

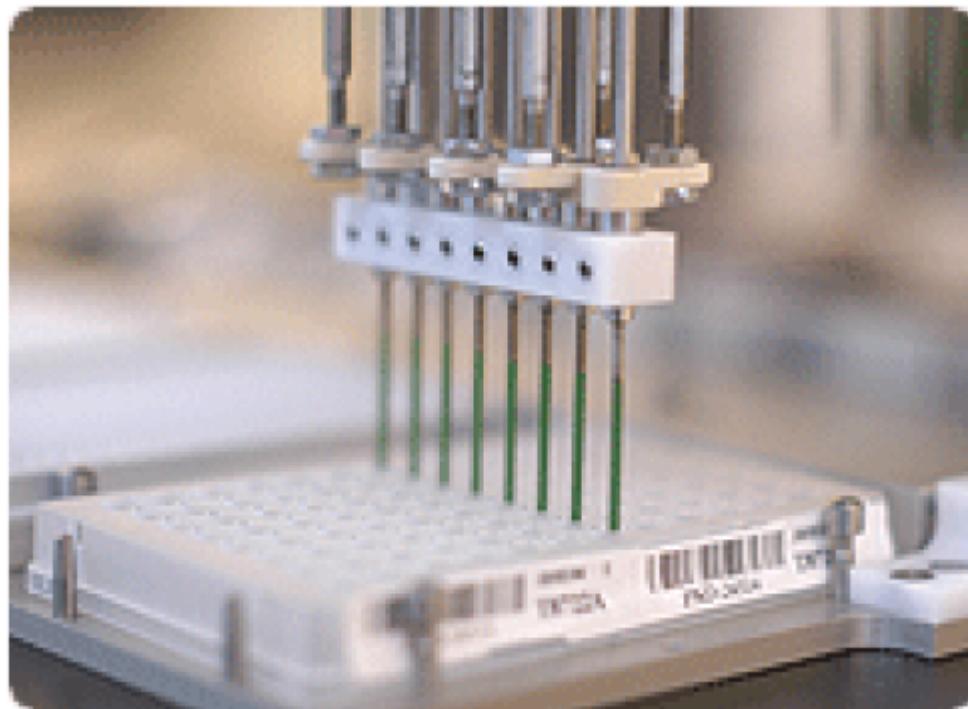
# Radio Link Frequency Assignment (CELAR)



(Cabon *et al*, Constraints 1999 ; CP97 – AAAI06 – IJCAI07 – IJCAI09 – CP10)

$n \leq 458$ ,  $d=44$ ,  $e(2) \leq 5,000$

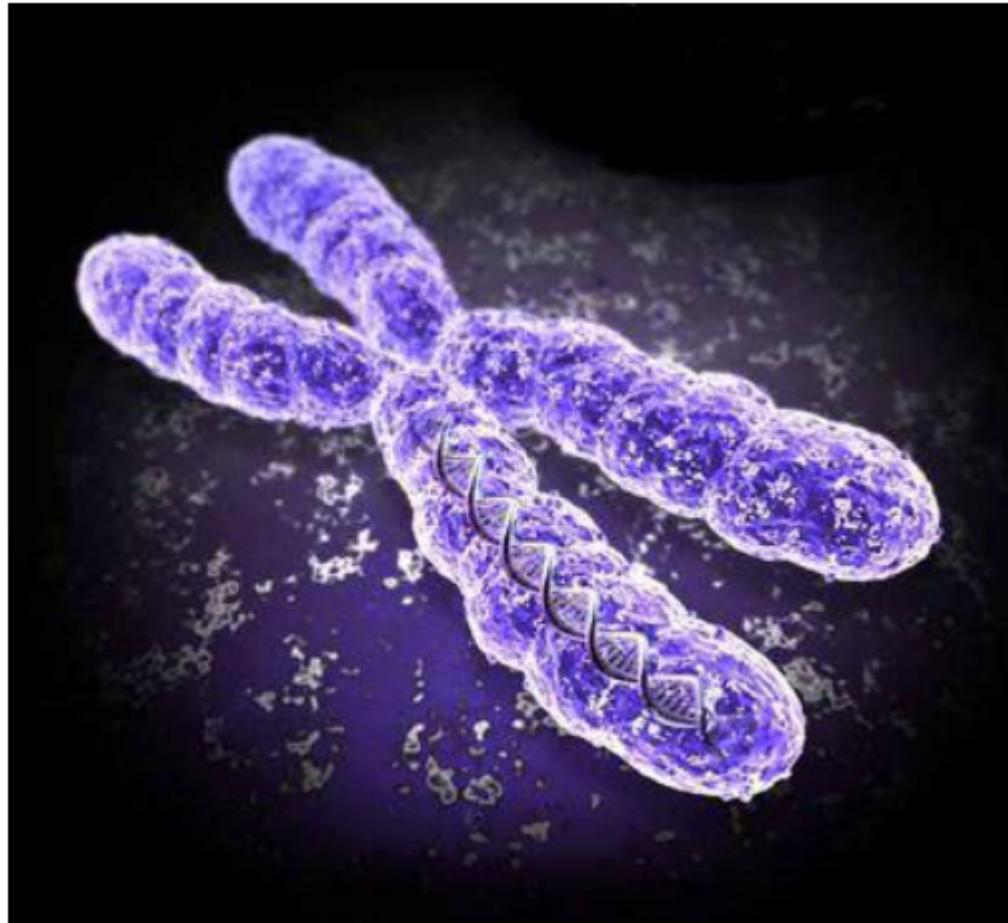
# Mendelian error correction in complex pedigree (MendelSoft)



(Constraints08)

$n \leq 20,000$ ,  $d \leq 66$ ,  $e(3) \leq 30,000$

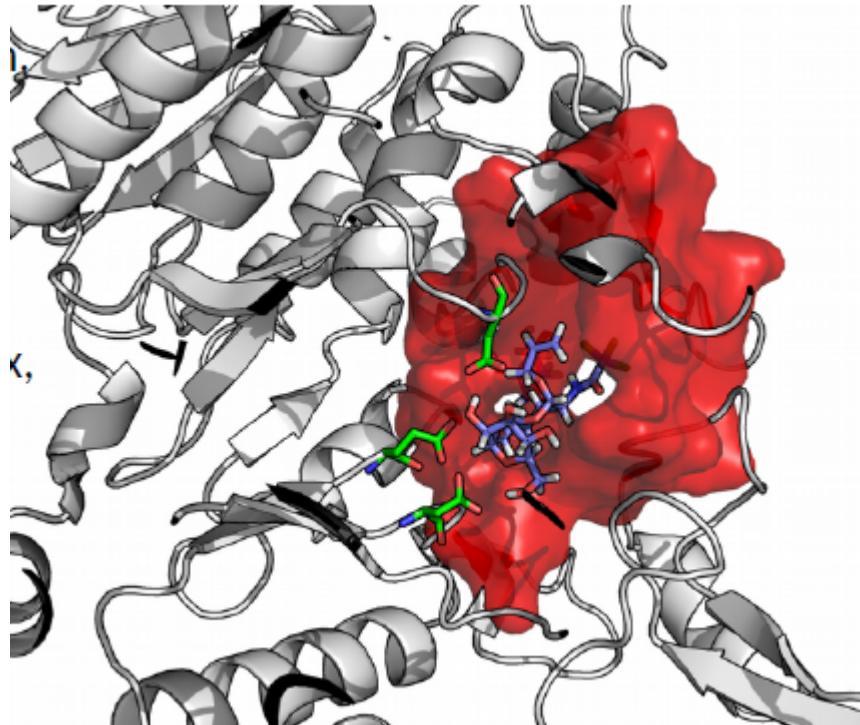
# Genetic Linkage Analysis



(Marinescu & Dechter, AAAI 2006 ; IJCAI11)

$n \leq 1,200$ ,  $d \leq 7$ ,  $e(2-5) \leq 2,000$

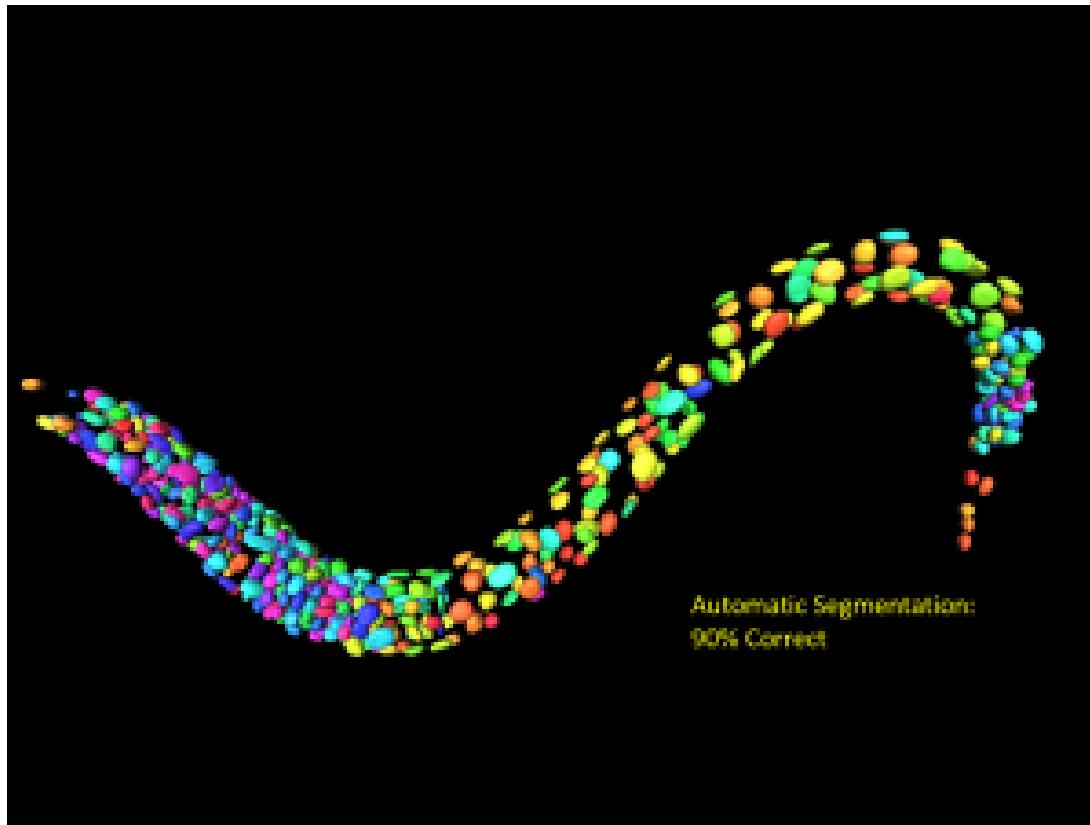
# Protein Design



(CP12 – Bioinformatics13 - **AIJ14** – JCTC15 – ISMP18)

$n \leq 120$ ,  $d \leq 190$ ,  $e(2) \leq 7,260$

# Graph Matching (worms segmentation)



(Kainmueller et al, Med Image Comput 2014 ; Haller et al, AAAI 2018)

$n \leq 558$ ,  $d \leq 128$ ,  $e(2) \leq 23,407$

# Graphical Model

## Definition (Graphical model)

- Let  $X = (X_1, \dots, X_n)$  be a set of variables.
- $X_i$  takes values in  $\Lambda_i \subseteq \mathbb{R}$ .
- A realization of  $X$  is denoted  $x = (x_1, \dots, x_n)$ , with  $x_i \in \Lambda_i$ .
- A graphical model over  $X$  is a function  $\psi : \prod_i \Lambda_i \rightarrow \mathbb{R}$ , which writes,  
 $\forall x \in X:$

$$\psi(x) = \odot_{B \in \mathcal{B}} \psi_B(x_B),$$

where  $\mathcal{B}$  is a set of subsets of  $V = \{1, \dots, n\}$ ,  $\psi_B : \prod_{i \in B} \Lambda_i \rightarrow \mathbb{R}$  and  $\odot \in \{\prod, \sum, \min, \max \dots\}$  is a combination operator.

# Probabilistic Graphical Models

## Definition (Markov chain)

- $X = (X_1, \dots, X_n)$  is a set of variables, with finite domains  $\{\Lambda_i\}_{i=1,\dots,n}$ .

$$P(x_1, \dots, x_n) = \underbrace{P(x_1)}_{\psi_1(x_1)} \times \underbrace{P(x_2|x_1)}_{\psi_{12}(x_1, x_2)} \times \dots \times \underbrace{P(x_n|x_{n-1})}_{\psi_{(n-1)n}(x_{n-1}, x_n)}$$

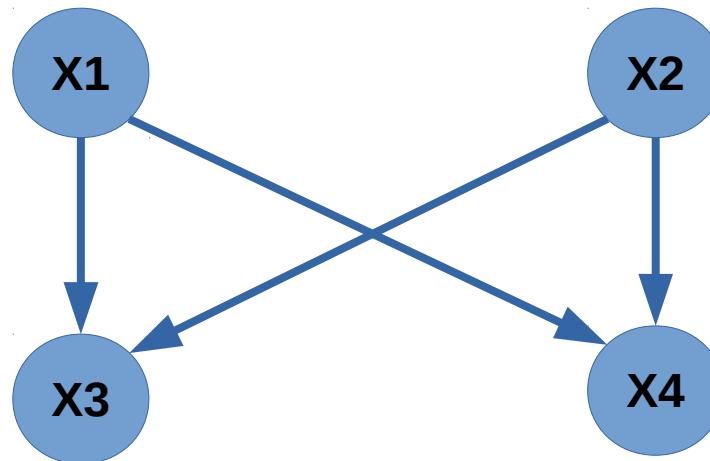


# Probabilistic Graphical Models

## Definition (Bayesian network)

- $X = (X_1, \dots, X_n)$  is a set of variables, with finite domains  $\{\Lambda_i\}_{i=1,\dots,n}$ .
- $Par(i) \subseteq \{1, \dots, i-1\}, \forall i = 2, \dots, n.$

$$P(x_1, \dots, x_n) = \underbrace{P(x_1)}_{\psi_1(x_1)} \times \prod_{i=2}^n \underbrace{P(x_i | x_{Par(i)})}_{\psi_{Par(i) \cup \{i\}}(x_i, x_{Par(i)})}$$

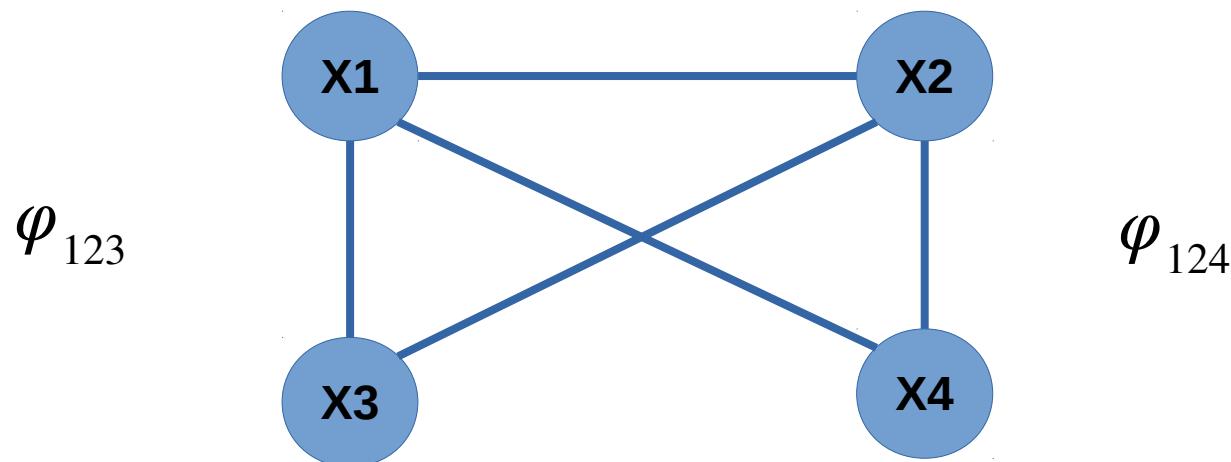


# Probabilistic Graphical Models

## Definition (Markov Random Field)

- $G = (V, E)$  is an undirected graph with vertices  $V = \{1, \dots, n\}$ , edges  $E \in V \times V$  and  $\mathcal{C}$  is the set of *cliques* of  $G$ .
- $\{\psi_C : X_C \rightarrow \mathbb{R}^{+*}\}_{C \in \mathcal{C}}$  are strictly positive functions.

$$P(x_1, \dots, x_n) = \underbrace{\frac{1}{Z}}_{\psi_\emptyset, \text{ normalizing constant}} \times \prod_{C \in \mathcal{C}} \psi_C(x_C)$$

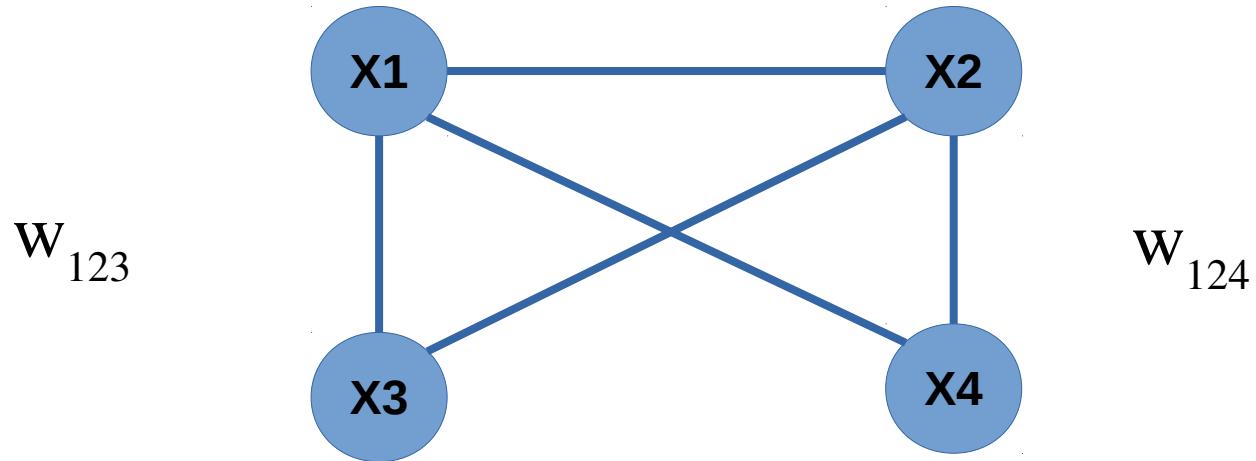


# Deterministic Graphical Model

## Definition (Cost Functions networks)

- $\{w_C : X_C \rightarrow \mathbb{R}^+\}_{C \in \mathcal{C}}$  are positive functions.

$$w(x_1, \dots, x_n) = \sum_{c \in \mathcal{C}} w_c(x_c)$$



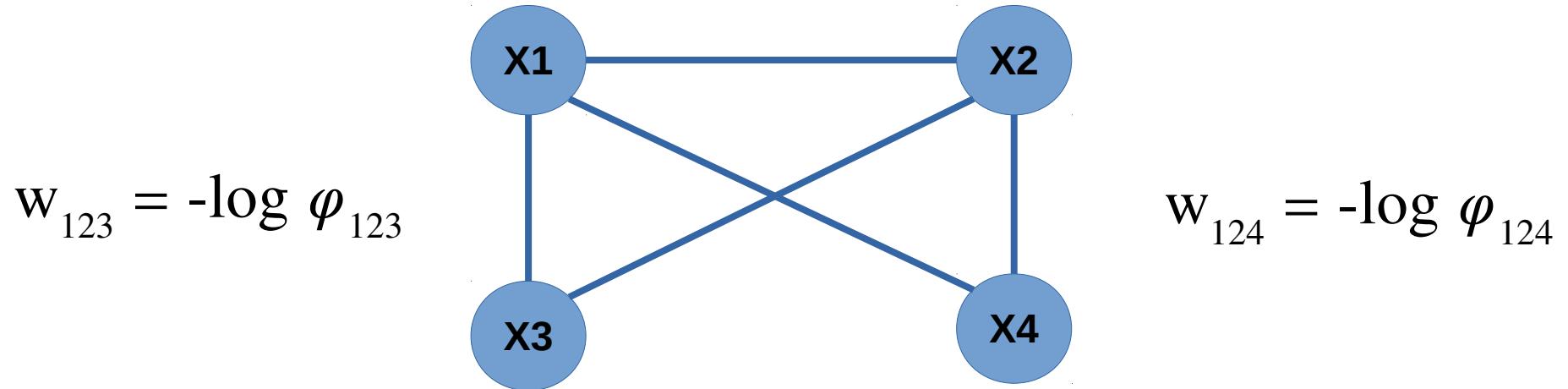
# Deterministic Graphical Model

Definition (Cost Functions networks)

- $\{w_C : X_C \rightarrow \mathbb{R}^+\}_{C \in \mathcal{C}}$  are positive functions.

$$w(x_1, \dots, x_n) = \sum_{c \in \mathcal{C}} w_c(x_c)$$

Minimization task:  $\min w(X_1, \dots, X_n)$       NP-hard problem



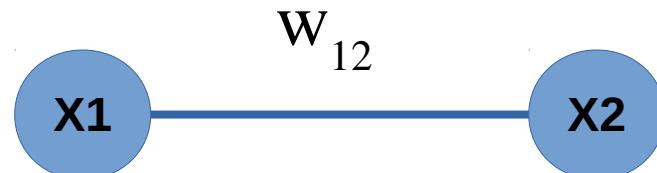
Energy minimization task is equivalent to finding the most probable explanation

# Example

In JSON compatible toolbar2 *cfn* format



```
{  
    problem: { name: "maximization", mustbe: ">-5.0"},  
    variables: { "X1": ["a", "b"], "X2": ["c", "d"] },  
    functions: {  
        "w0": {scope: [], costs: [-6.0]},  
        "w1": {scope: ["X1"], costs: [1.0, 0.5]},  
        "w2": {scope: ["X2"], costs: [1.0, 0.5]},  
        "w12": {scope: ["X1", "X2"], costs: [-1.0, 0.5, -2.0, 5.5]}  
    }  
}
```

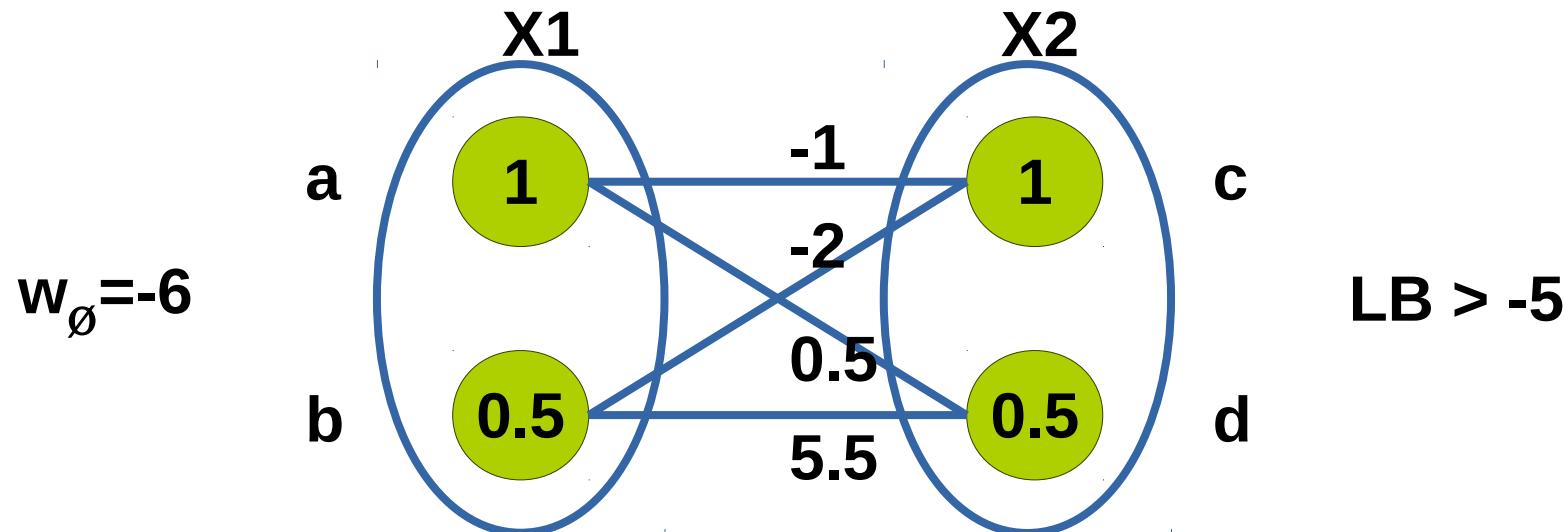


# Micro-Structure

{

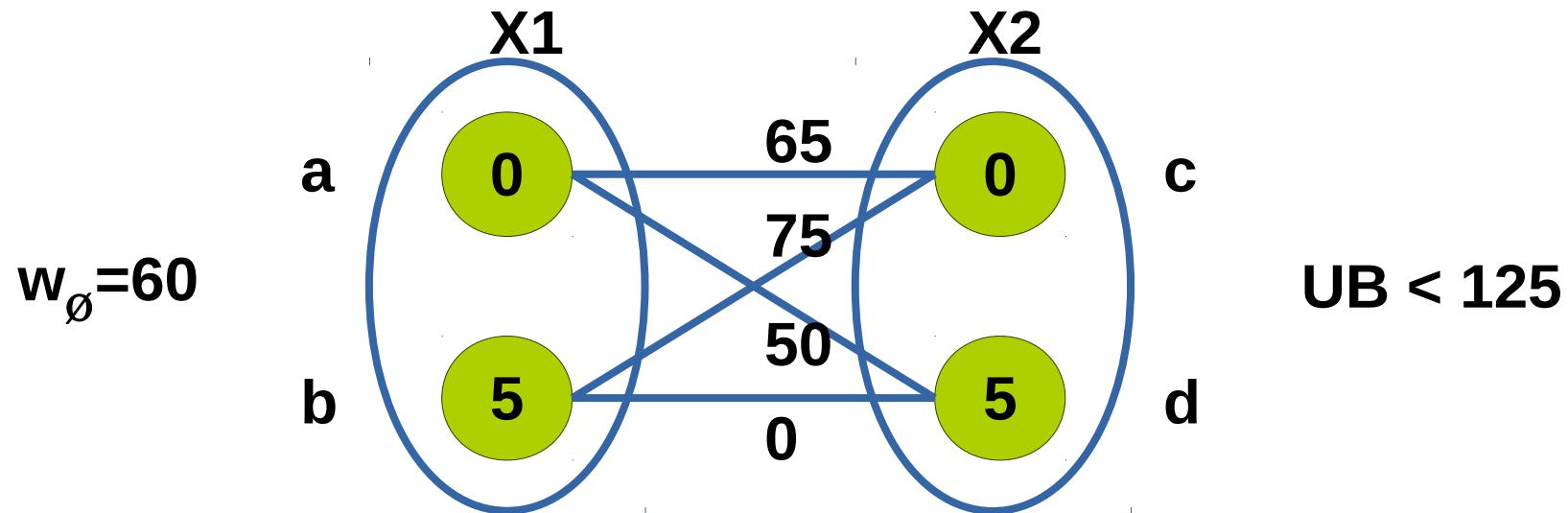
```
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    "w12": {scope: ["X1", "X2"], costs: [-1.0, 0.5, -2.0, 5.5]}  
}
```

}



# Minimization with non-negative integer costs

```
{  
    problem: { name: "maximization", mustbe: ">-5.0"},  
    variables: { "X1": ["a", "b"], "X2": ["c", "d"] },  
    functions: {  
        "w0": {scope: [], costs: [-6.0]},  
        "w1": {scope: ["X1"], costs: [1.0, 0.5]},  
        "w2": {scope: ["X2"], costs: [1.0, 0.5]},  
        "w12": {scope: ["X1", "X2"], costs: [-1.0, 0.5, -2.0, 5.5]}  
    }  
}
```

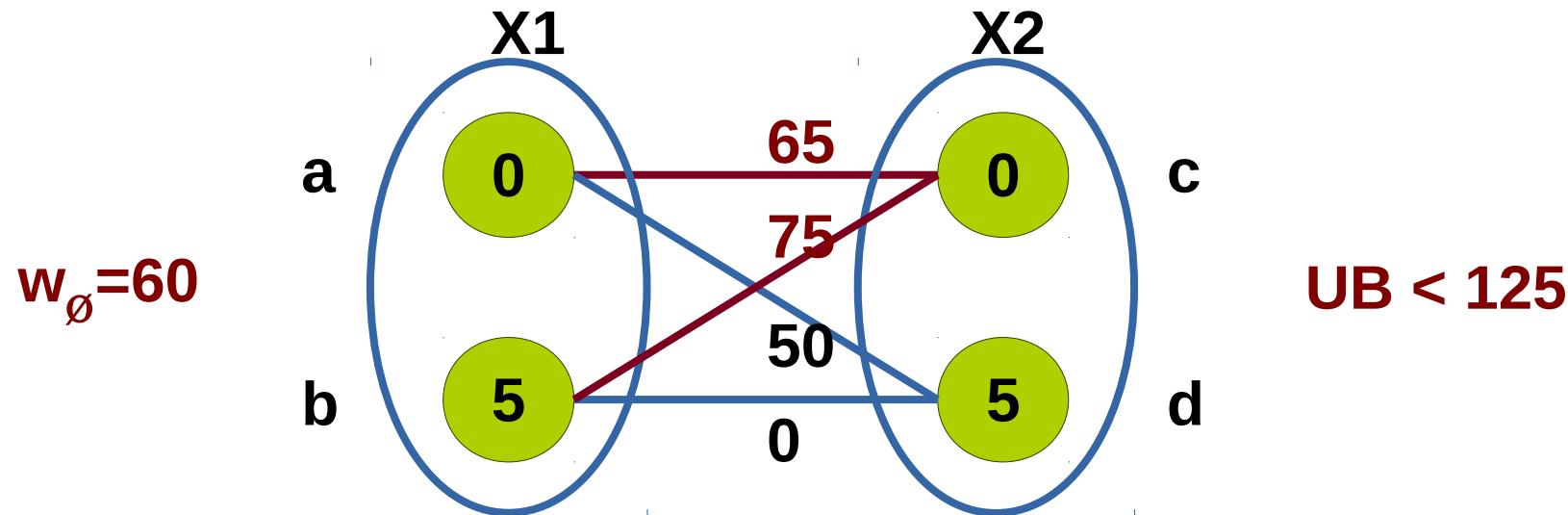


# Constraints are Cost Functions

{

```
problem: { name: "maximization", mustbe: ">-5.0"},  
variables: { "X1": ["a", "b"], "X2": ["c", "d"] },  
functions: {  
    "w0": {scope: [], costs: [-6.0]},  
    "w1": {scope: ["X1"], costs: [1.0, 0.5]},  
    "w2": {scope: ["X2"], costs: [1.0, 0.5]},  
    "w12": {scope: ["X1", "X2"], costs: [-1.0, 0.5, -2.0, 5.5]}  
}
```

}

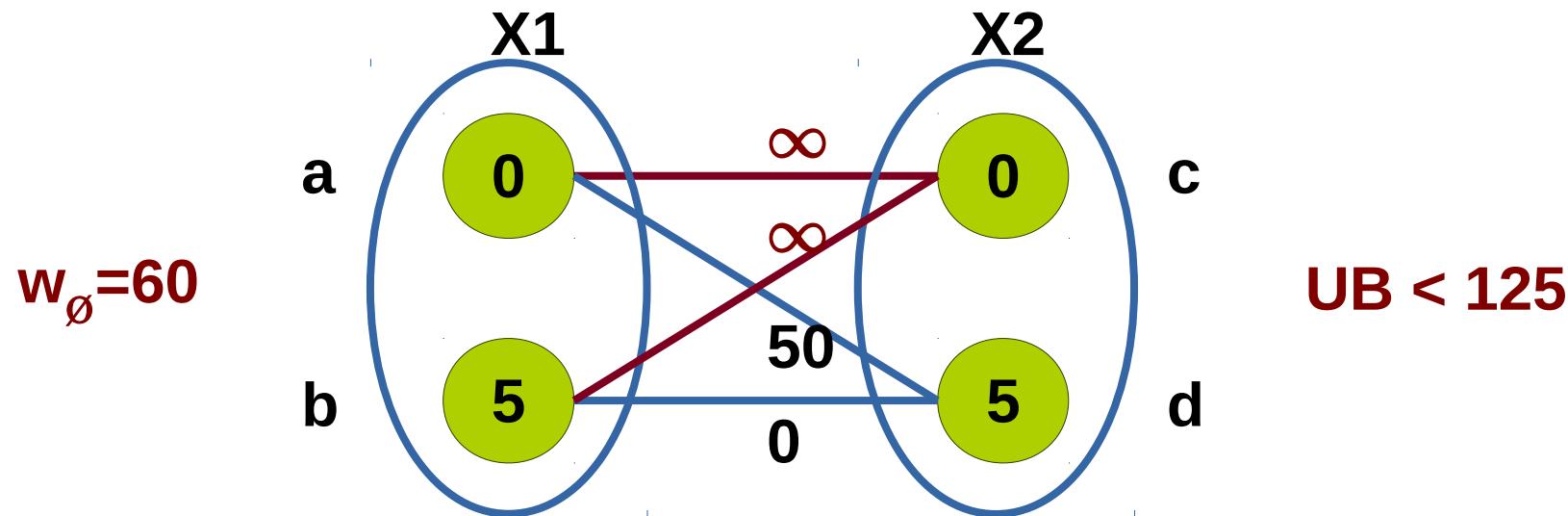


# Constraints are Cost Functions

{

```
problem: { name: "maximization", mustbe: ">-5.0"},  
variables: { "X1": ["a", "b"], "X2": ["c", "d"] },  
functions: {  
    "w0": {scope: [], costs: [-6.0]},  
    "w1": {scope: ["X1"], costs: [1.0, 0.5]},  
    "w2": {scope: ["X2"], costs: [1.0, 0.5]},  
    "w12": {scope: ["X1", "X2"], costs: [-1.0, 0.5, -2.0, 5.5]}  
}
```

}



# Other equivalent formulations

*In various toulbar2 input formats*

- WCSP

```
wcsp 2 2 4 125  
2 2  
2 0 1 0 4  
0 0 65  
0 1 50  
1 0 75  
1 1 0  
1 0 125 2  
0 0  
1 5  
1 1 125 2  
0 0  
1 5  
0 6 0 0
```

- MRF

```
MARKOV  
2  
2 2  
4  
2 0 1  
1 0  
1 1  
1 0  
4  
0.000341454887383  
0.00215443469003  
0.0001  
1.0  
2  
1.0  
0.541169526546  
2  
1.0  
0.541169526546  
2  
0.00063095734448  
0.00063095734448
```

- Max-SAT

```
p wcnf 2 7 125  
65 1 2 0  
50 1 -2 0  
75 -1 2 0  
5 -1 0  
5 -2 0  
60 1 0  
60 -1 0
```

- QPBO

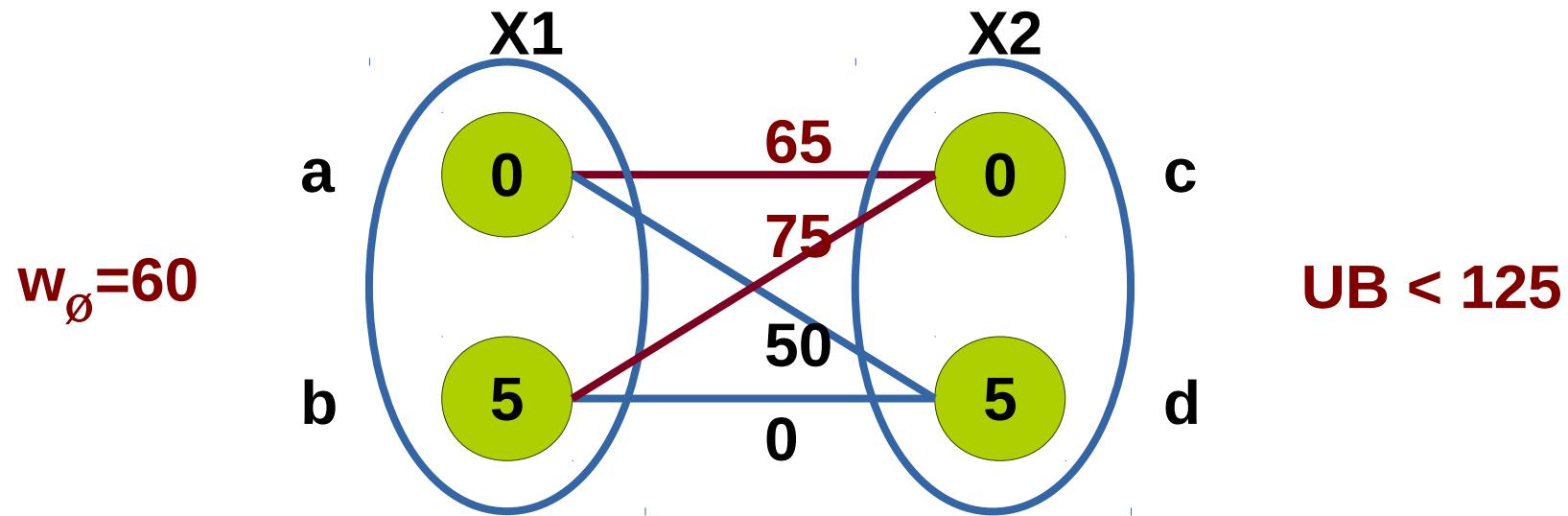
```
4 13  
1 3 32.5  
1 4 25  
2 3 37.5  
2 2 5  
4 4 5  
1 1 60  
2 2 60  
1 1 -1000  
2 2 -1000  
1 2 1000  
3 3 -1000  
4 4 -1000  
3 4 1000
```

# Local reasoning techniques

*Cost Function Propagation*

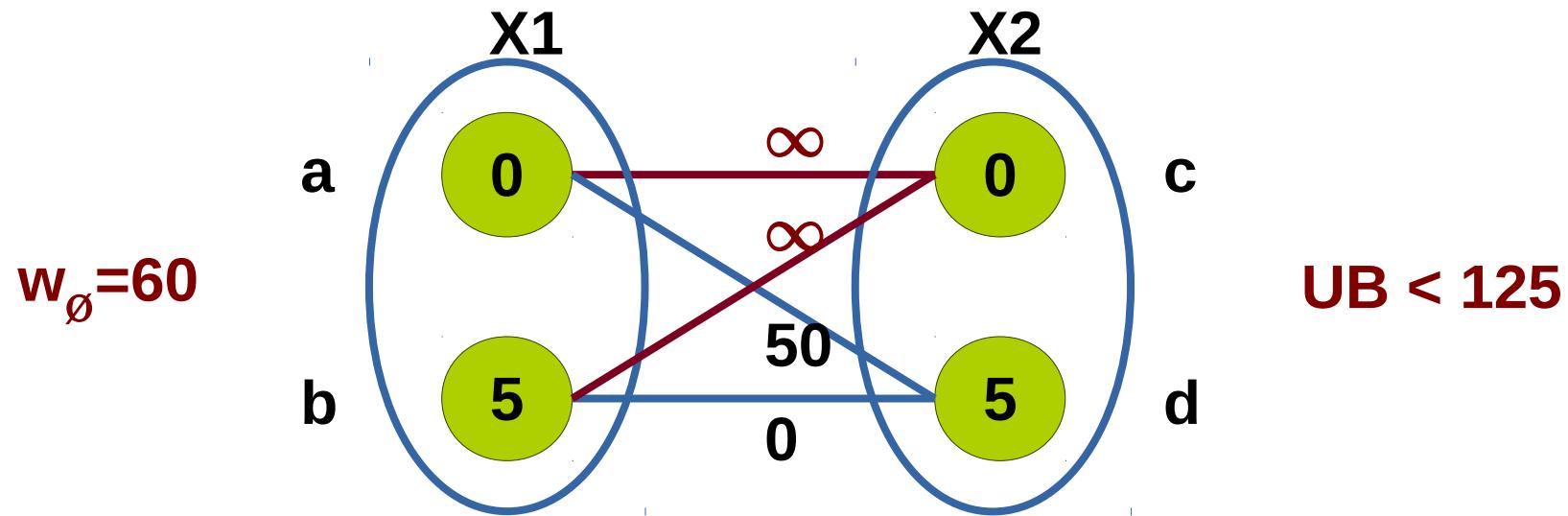
# Reparameterization and pruning

(Schiex, CP 2000 ; Larrosa, AAAI 2002 ; Cooper, FSS 2003 ; IJCAI05 ; IJCAI07 ; AAAI08 ; AIJ10)



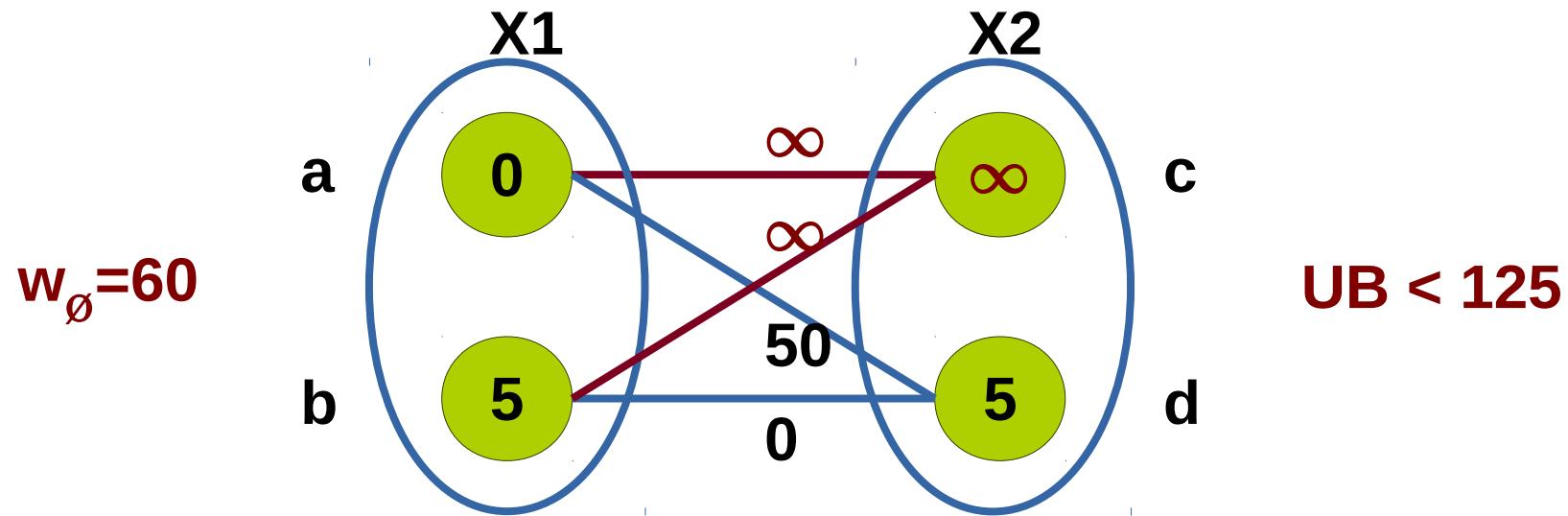
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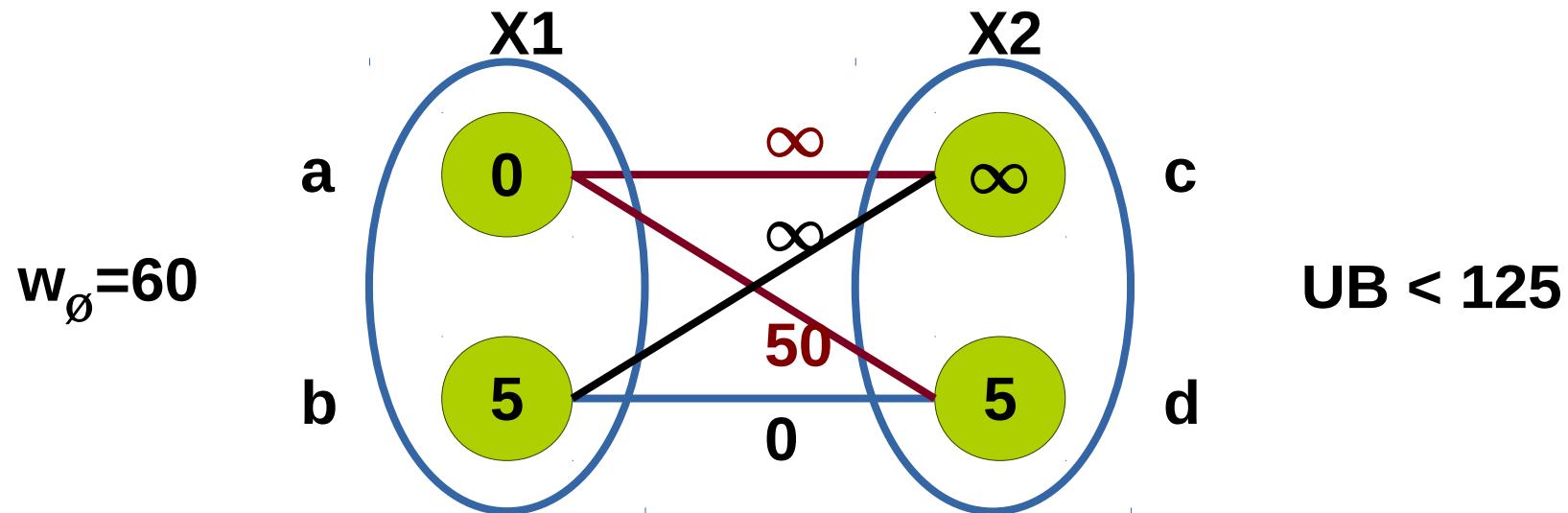
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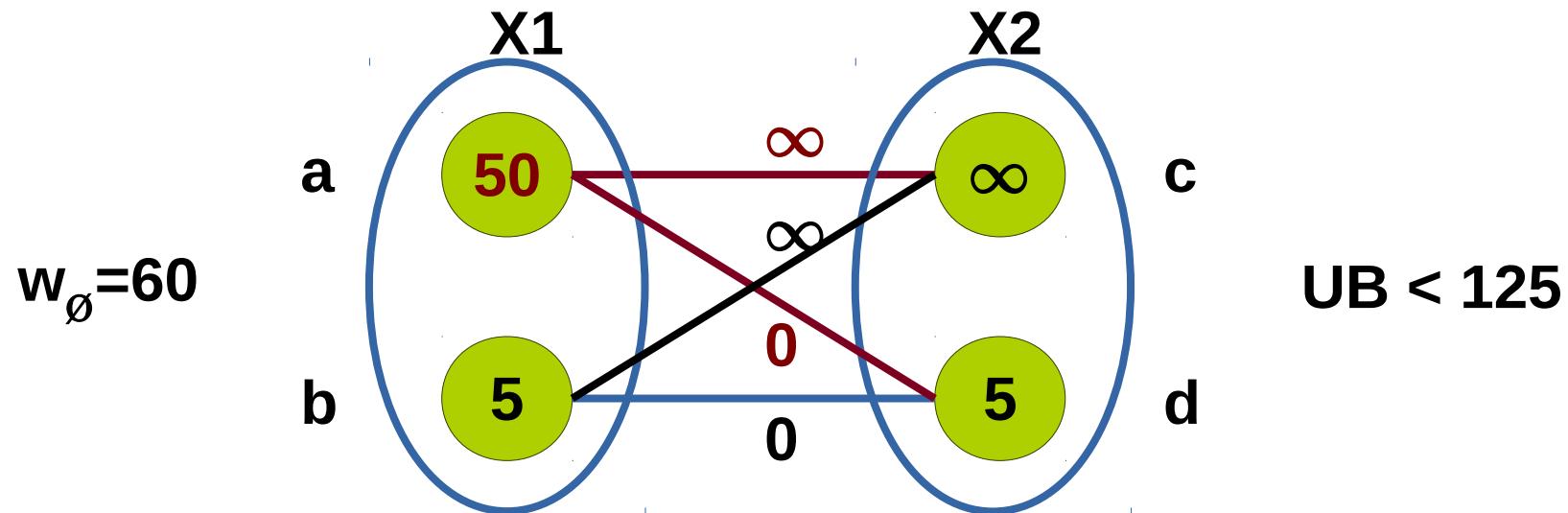
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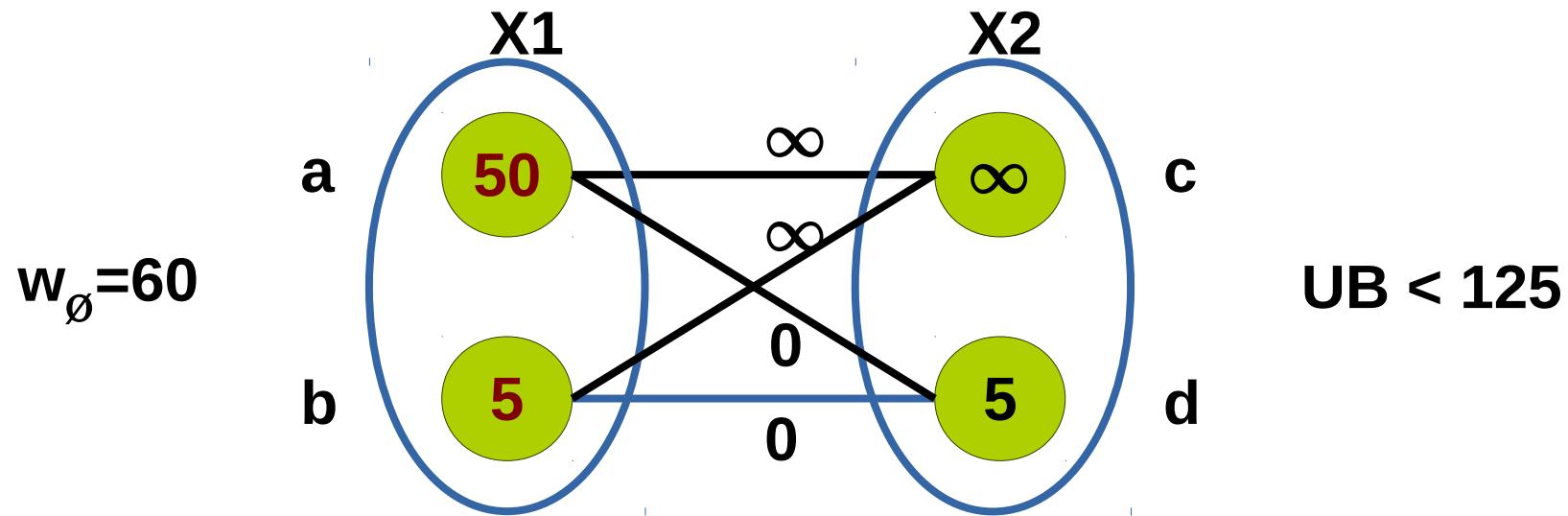
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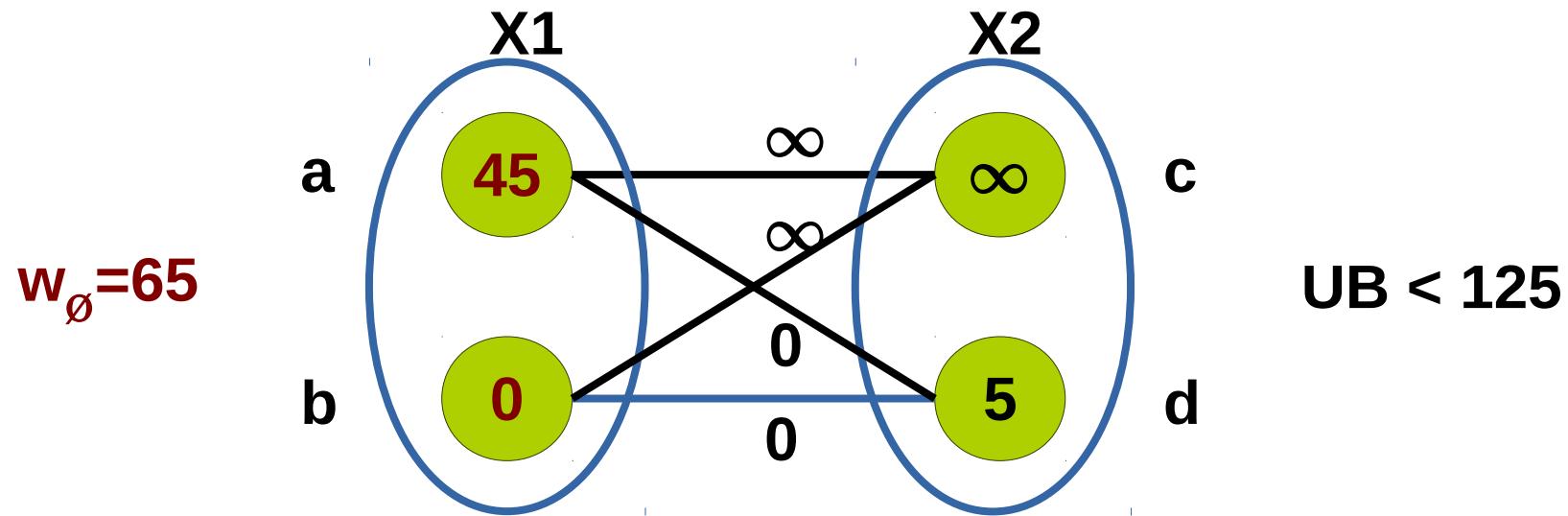
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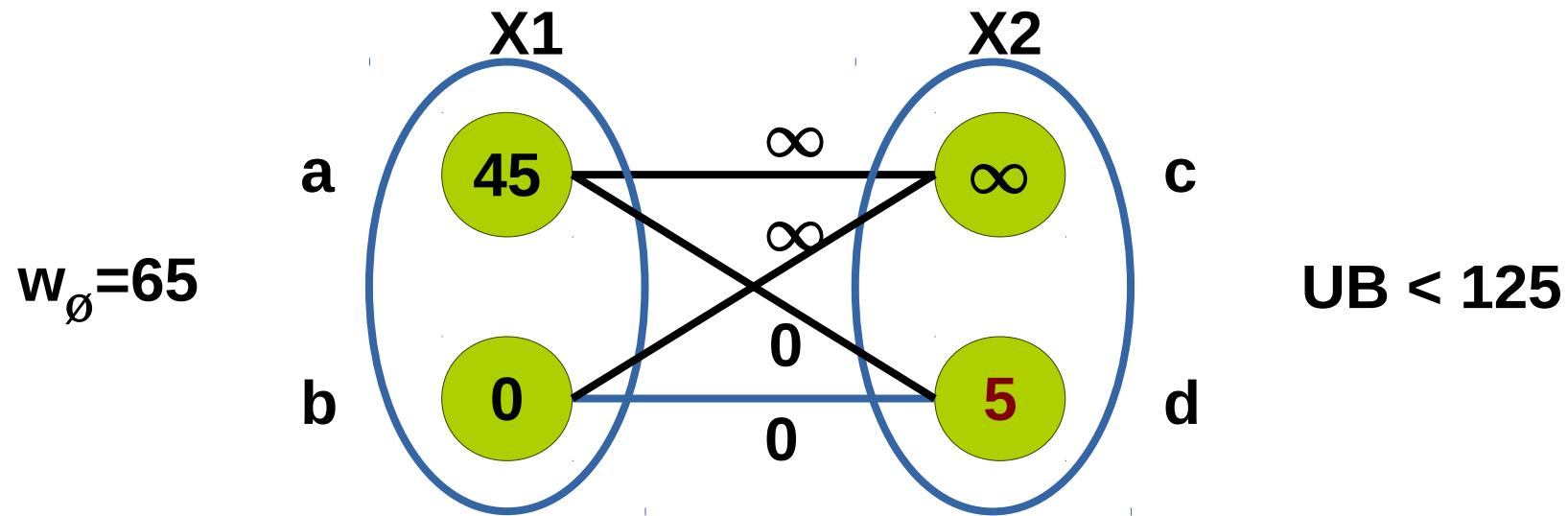
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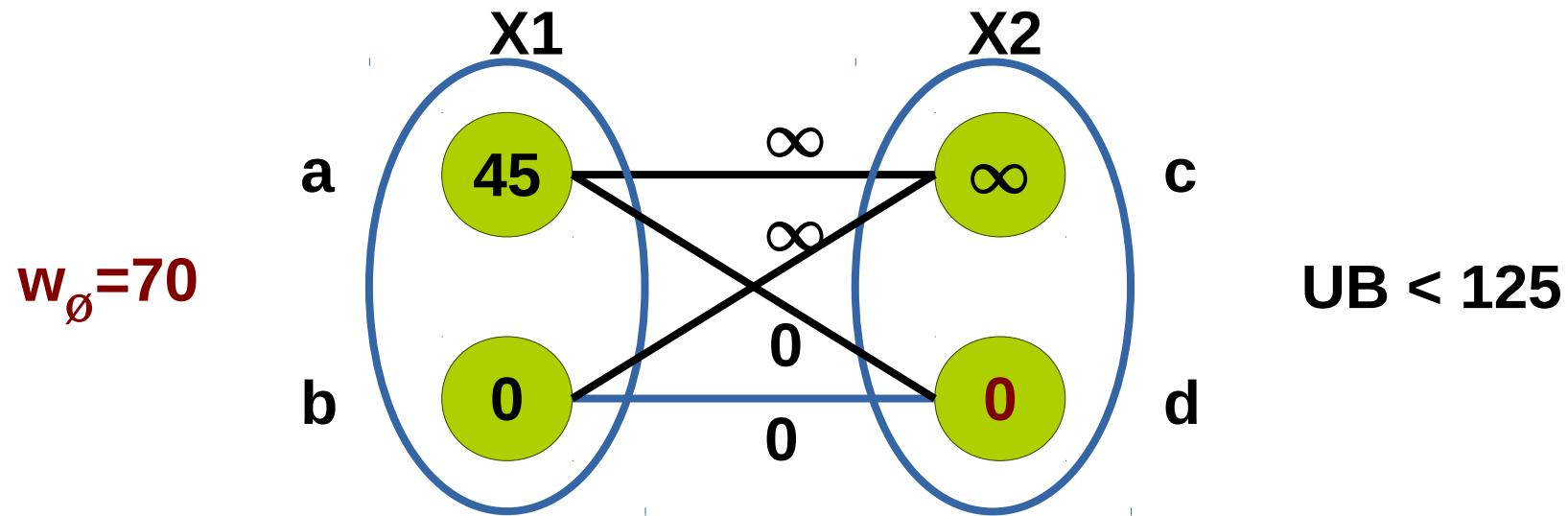
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(Schiex, CP 2000 ; Larrosa, AAAI 2002 ; Cooper, FSS 2003 ; IJCAI05 ; IJCAI07 ; AAAI08 ; AIJ10)



# Reparameterization and pruning

(Schiex, CP 2000 ; Larrosa, AAAI 2002 ; Cooper, FSS 2003 ; IJCAI05 ; IJCAI07 ; AAAI08 ; AIJ10)



# Reparameterization and pruning

(Schiex, CP 2000 ; Larrosa, AAAI 2002 ; Cooper, FSS 2003 ; IJCAI05 ; IJCAI07 ; AAAI08 ; AIJ10)

- Reparameterization produces a feasible solution of the **dual** of a strong LP relaxation
  - We use a **sequence** of reparameterizations
    - Faster than LP
    - Not optimal: weaker dual bounds than LP
    - Many fixpoints
- and domain value pruning

# Same Example in 01LP

(CPAIOR16 – Constraints16)

- **Direct LP formulation**

Minimize

$$+50 t_{0\_0\_1\_1} + 75 t_{0\_1\_1\_0} + 65 t_{0\_0\_1\_0} - 5 d_{0\_0} - 5 d_{1\_0} + 60 t + 10 t$$

Subject to:

$$+1 d_{0\_0} - 1 d_{1\_0} - t_{0\_0\_1\_1} \leq 0$$

$$-1 d_{0\_0} + 1 d_{1\_0} - t_{0\_1\_1\_0} \leq 0$$

$$+1 d_{0\_0} + 1 d_{1\_0} - t_{0\_0\_1\_0} \leq 1$$

Bounds

$$t_{0\_0\_1\_0} \leq 1$$

$$t_{0\_0\_1\_1} \leq 1$$

$$t_{0\_1\_1\_0} \leq 1$$

$$t = 1$$

Binary

$$d_{0\_0} \quad d_{1\_0}$$

End

- **Stronger LP formulation**

Minimize

$$+50 t_{0\_0\_1\_1} + 75 t_{0\_1\_1\_0} + 65 t_{0\_0\_1\_0} - 5 d_{0\_0} - 5 d_{1\_0} + 60 t + 10 t$$

Subject to:

$$+1 t_{0\_0\_1\_0} + 1 t_{0\_0\_1\_1} - 1 d_{0\_0} = 0$$

$$+1 t_{0\_1\_1\_0} + 1 t_{0\_1\_1\_1} + 1 d_{0\_0} = 1$$

$$+1 t_{0\_0\_1\_0} + 1 t_{0\_1\_1\_0} - 1 d_{1\_0} = 0$$

$$+1 t_{0\_0\_1\_1} + 1 t_{0\_1\_1\_1} + 1 d_{1\_0} = 1$$

Bounds

$$t_{0\_0\_1\_0} \leq 1$$

$$t_{0\_0\_1\_1} \leq 1$$

$$t_{0\_1\_1\_0} \leq 1$$

$$t_{0\_1\_1\_1} \leq 1$$

$$t = 1$$

Binary

$$d_{0\_0} \quad d_{1\_0}$$

End

# Uncapacitated Warehouse Location Problem

(Kratica et al., RAIRO OR 2001)

## Search nodes

Instance	cplex 12.7.1	toulbar2 1.0.0
Capmo1 100x100	155	7,581
Capmo2 100x100	25	2,024
Capmo3 100x100	93	5,439
Capmo4 100x100	23	4,055
Capmo5 100x100	28	2,664

## CPU time (sec. on PC i7 3GHz)

Instance	cplex 12.7.1	toulbar2 1.0.0
Capmo1 100x100	13.01	20.13
Capmo2 100x100	3.06	3.02
Capmo3 100x100	13.32	11.40
Capmo4 100x100	3.26	7.45
Capmo5 100x100	2.68	4.62

# Clique cuts

Given a set  $S$

$$x_i + x_j \leq 1 \quad \forall x_i, x_j \in S$$

$\Rightarrow$  Satisfied by  $x_i = 0.5$

But we can get

$$\sum_{x_i \in S} x_i \leq 1$$

# Clique cuts in CFN

(CP17)

Straightforward generalization

Given a set  $S$  of  $\langle X_i, v_i \rangle$  with

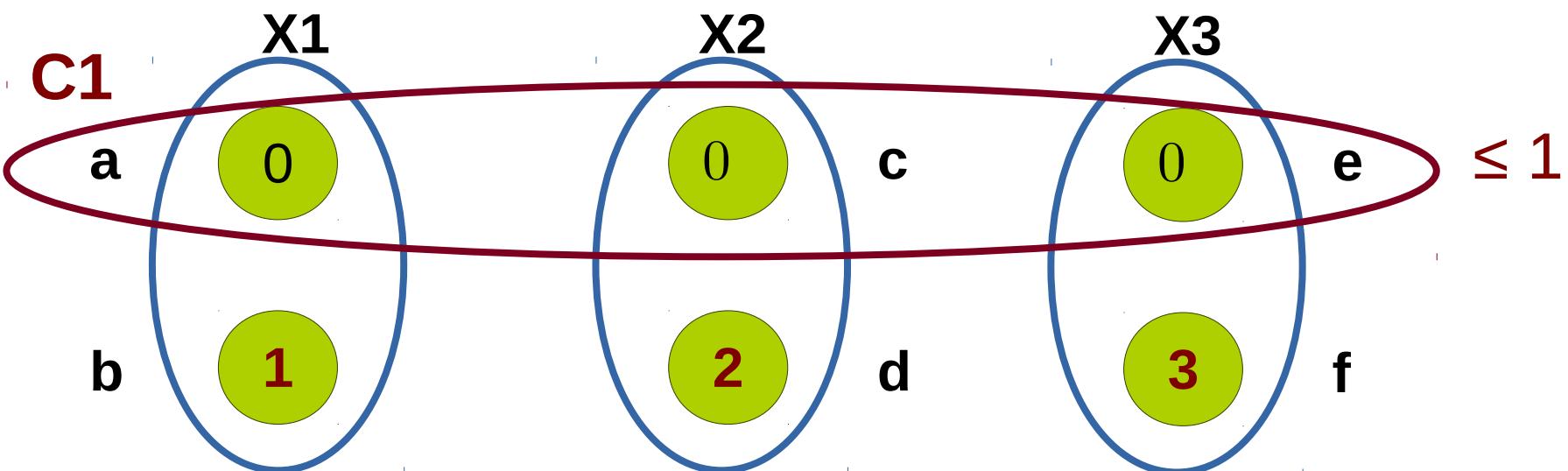
- $c_{ij}(v_i, v_j) = \infty$

Then derive

$$\sum_{ij \in S} x_{ij} \leq 1$$

# Reparameterization for clique

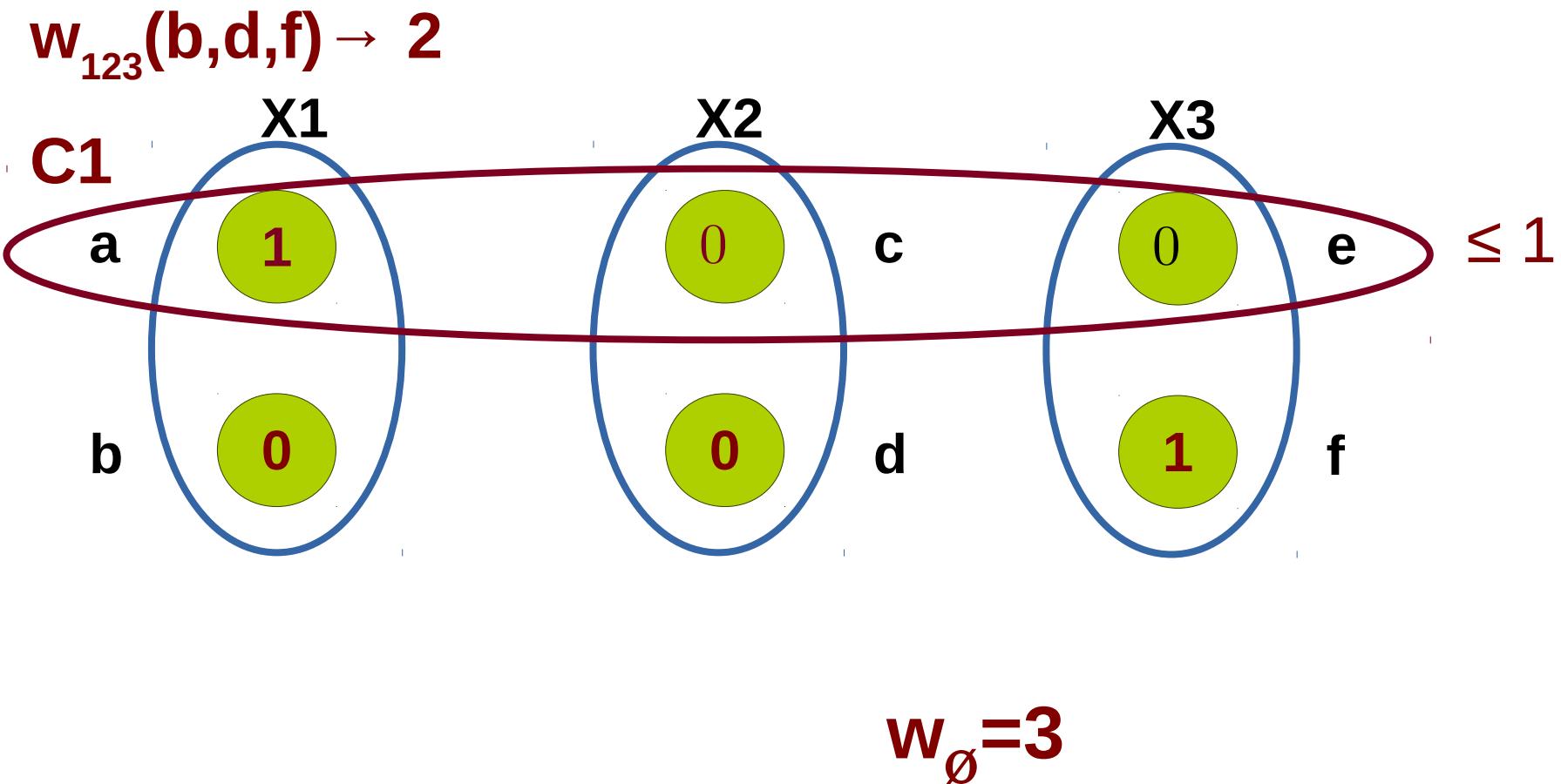
(CP17)



$$w_{\emptyset} = 0$$

# Reparameterization for clique

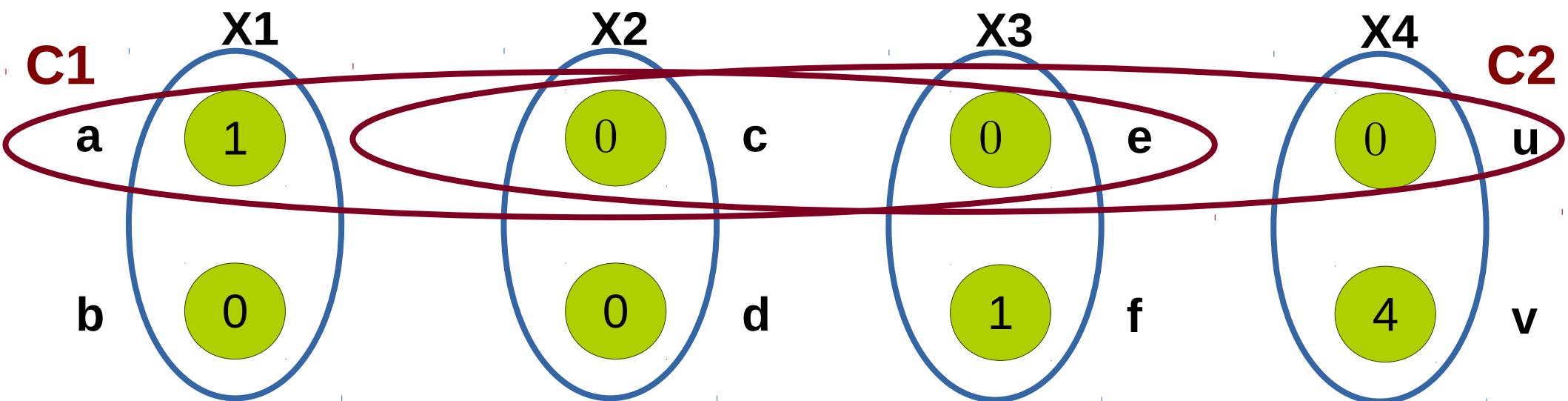
(CP17)



# Reparameterization for cliques

(CP17)

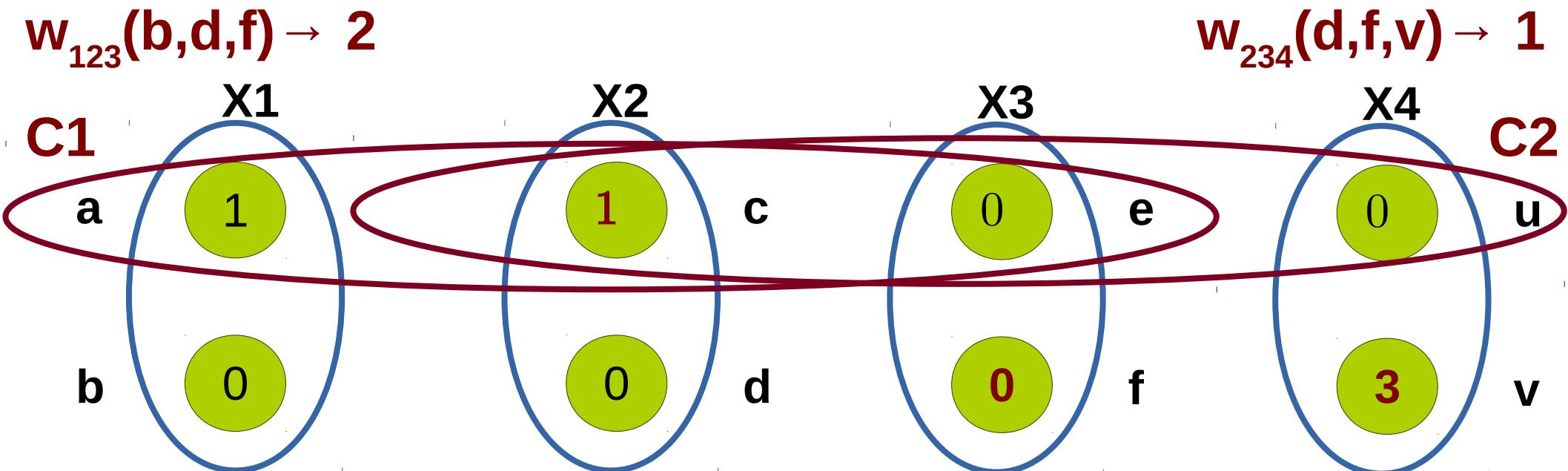
$$w_{123}(b,d,f) \rightarrow 2$$



$$w_\emptyset = 3$$

# Reparameterization for cliques

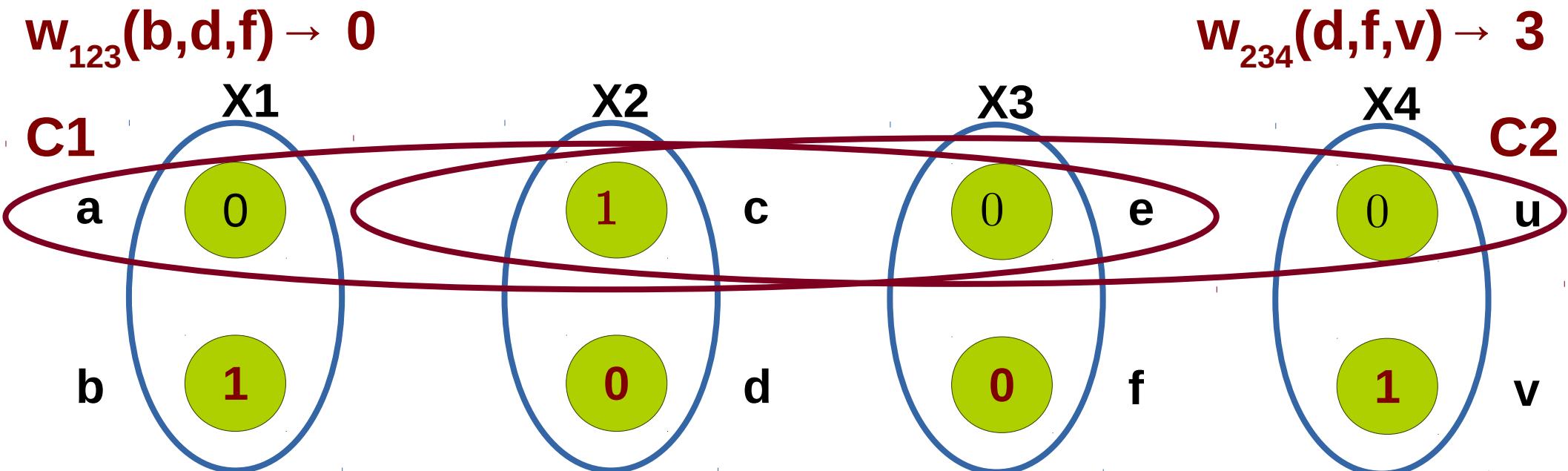
(CP17)



*Propagating  $C_1$  before  $C_2$*

# Reparameterization for cliques

(CP17)



# Experimental Results

(CP17)

problem	TOULBAR2		TOULBAR2 <sup>clq</sup>		CPLEX	
	solv.	time	solv.	time <sup>*</sup>	solv.	time
Auction/path	<b>86</b>	59	<b>86</b>	0.18	<b>86</b>	0.01
Auction/sched	<b>84</b>	110	<b>84</b>	0.23	<b>84</b>	0.04
MaxClique	31	1871	37	1508	<b>38</b>	<b>1533</b>
SPOT5	4	2884	6	2603	<b>16</b>	738

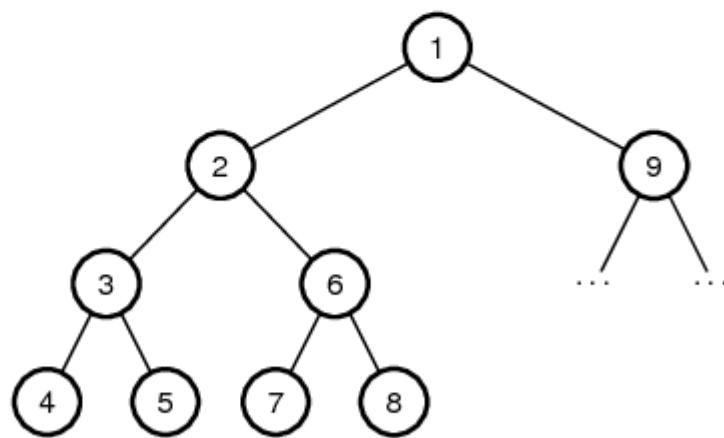
\* Including bounded clique detection with Bron-Kerbosch algorithm in preprocessing

# Complete tree search methods

*Hybrid search*

# DFS

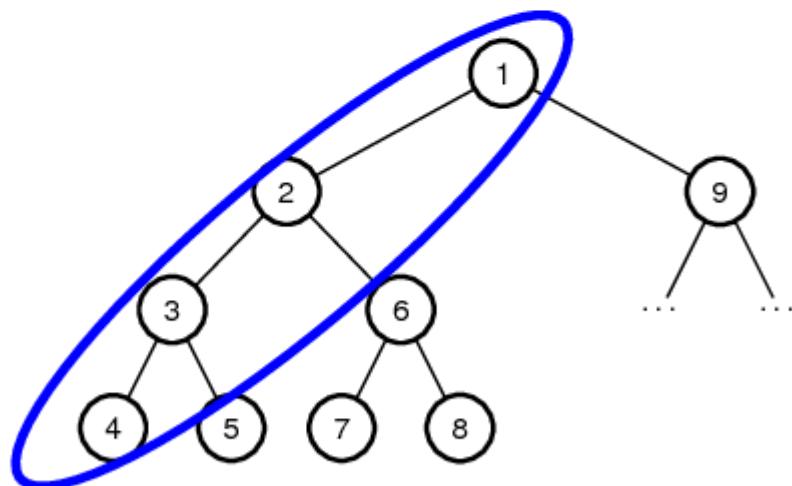
Depth First



# DFS

## Depth First Advantages

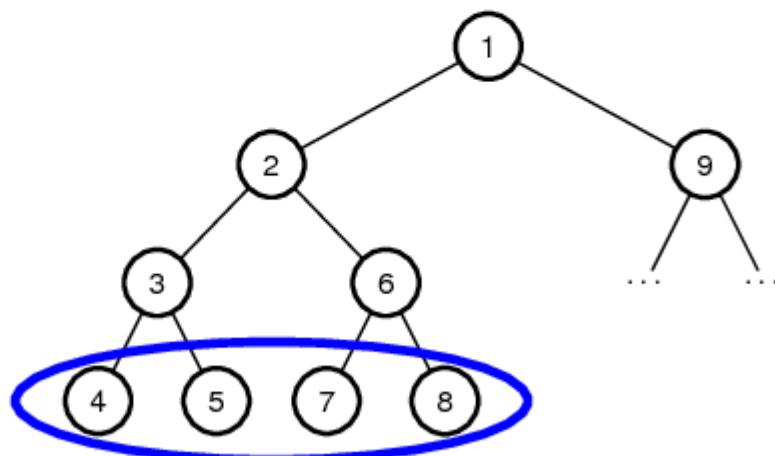
- Incrementality



# DFS

## Depth First Advantages

- Incrementality
- Anytime (sort of)



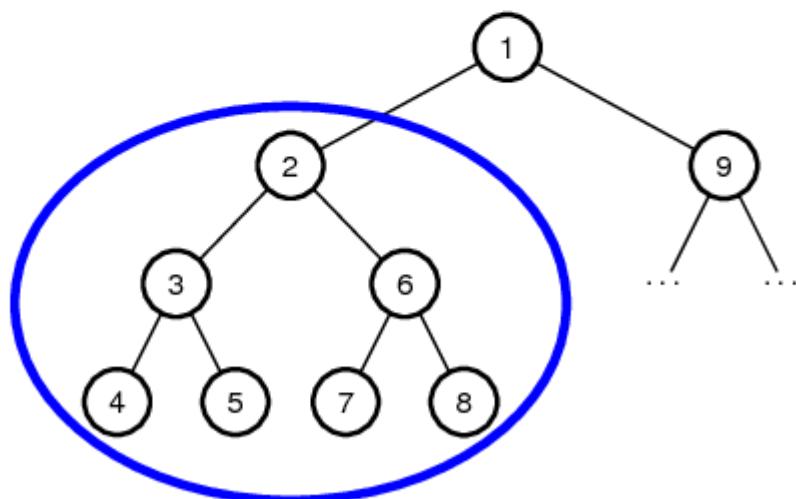
# DFS

## Depth First Advantages

- Incrementality
- Anytime (sort of)

But

- Thrashing



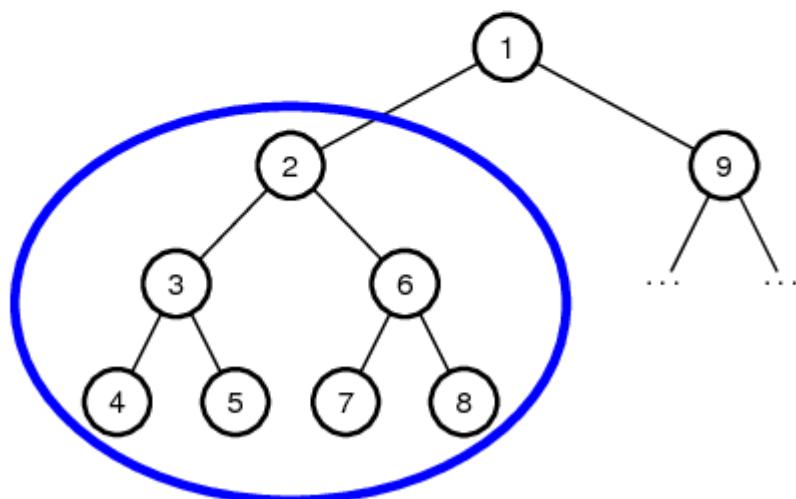
# DFS

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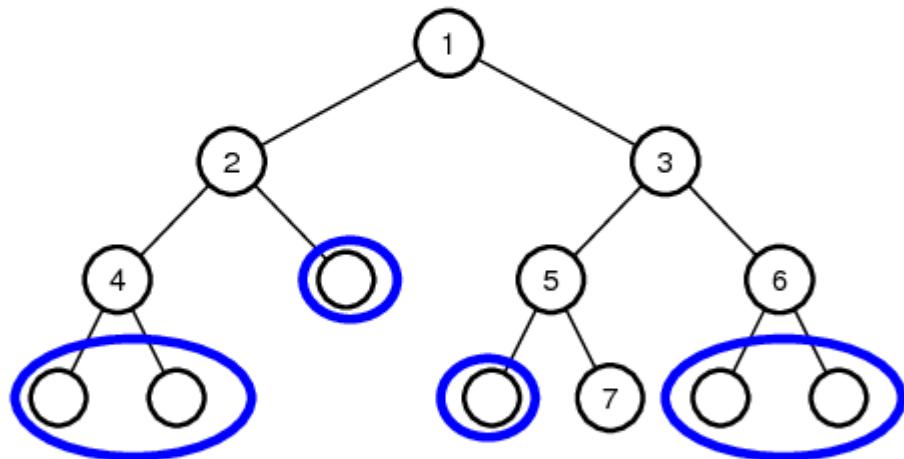
- Thrashing
- No global lower bounds



# BFS

Best first

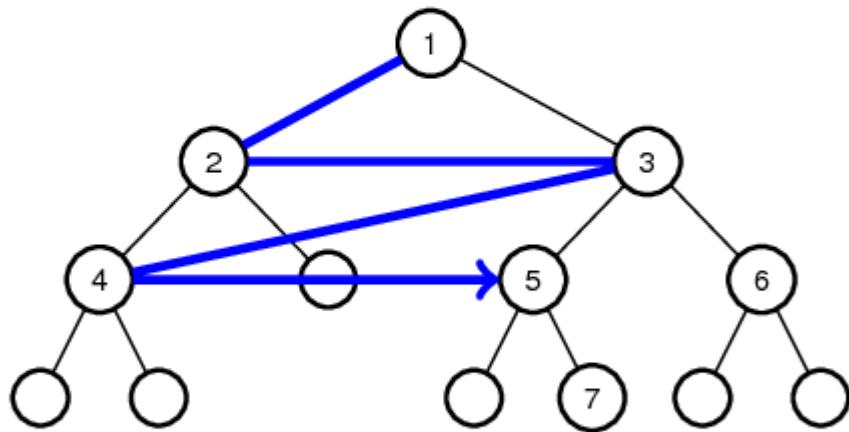
- Memory requirements



# BFS

Best first

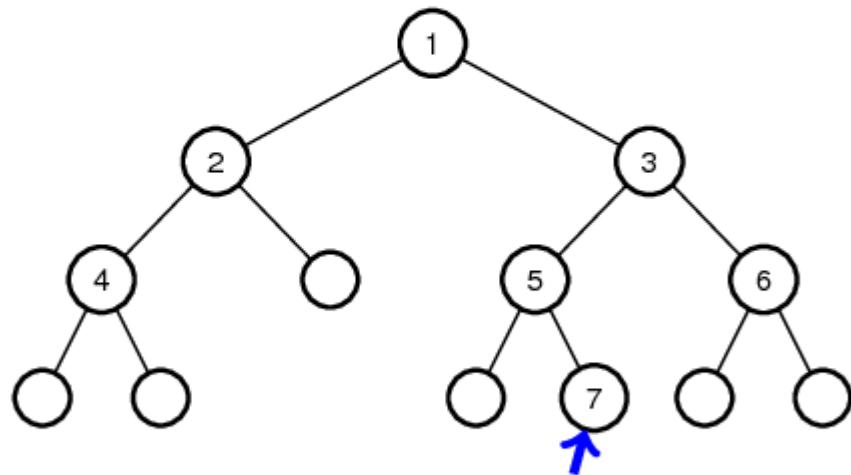
- Memory requirements
- No incrementality or even greater memory cost



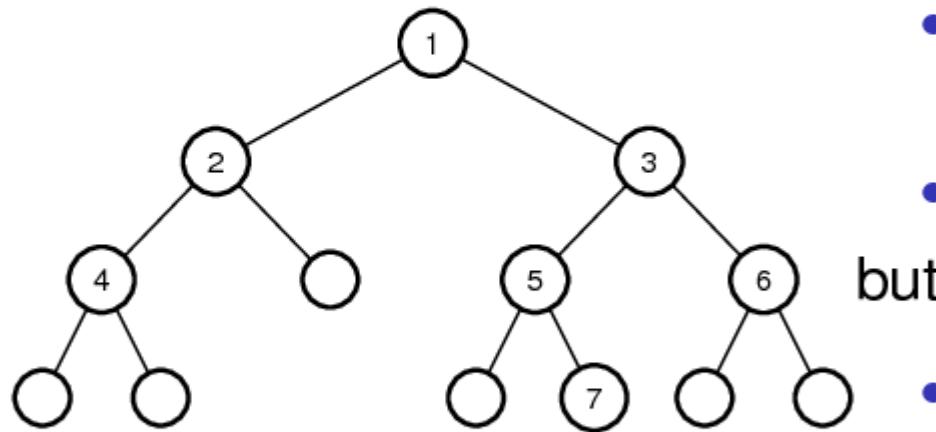
# BFS

Best first

- Memory requirements
- No incrementality or even greater memory cost
- Not anytime



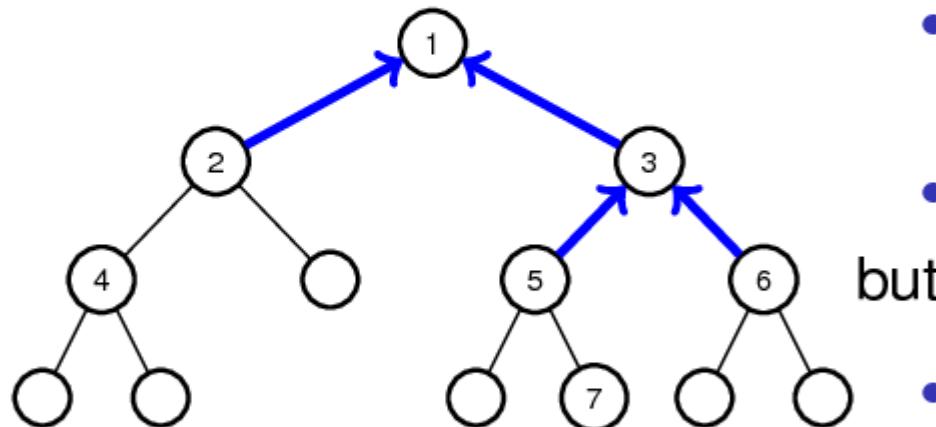
# BFS



Best first

- Memory requirements
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  - Not anytime
- but
- Theoretical guarantees

# BFS

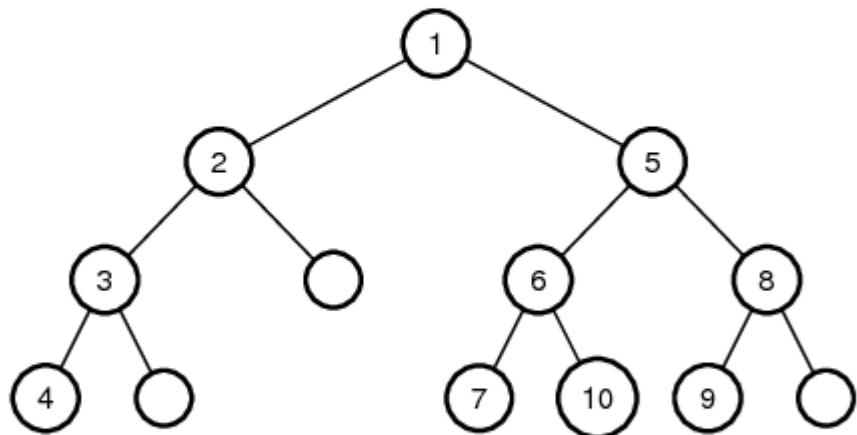


Best first

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  - No incrementality or even greater memory cost
  - Not anytime
- but
- Theoretical guarantees
  - Global lower bounds

# HBFS

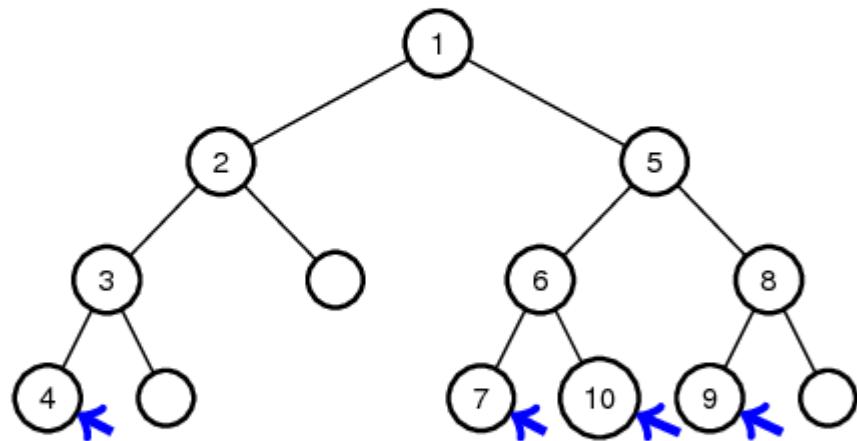
BFS with DFS probes\*



# HBFS

BFS with DFS probes\*

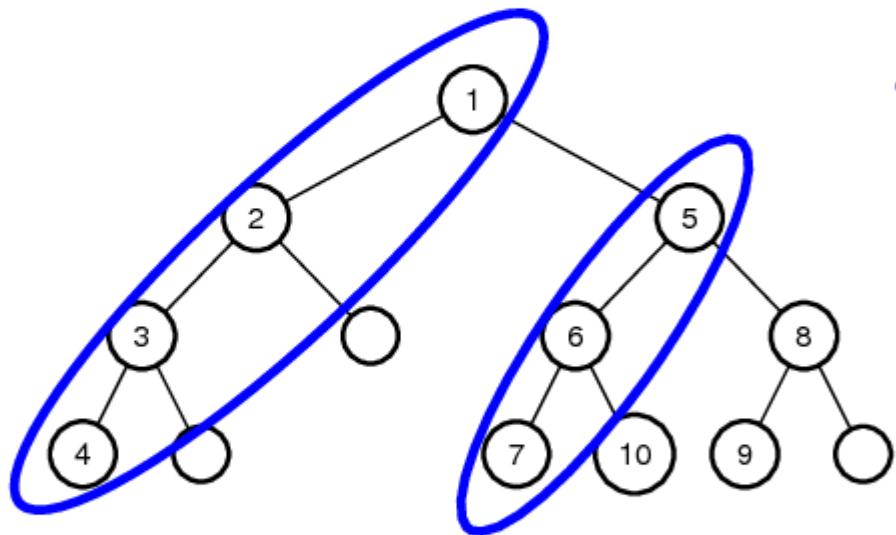
- Improved anytime behavior



# HBFS

BFS with DFS probes\*

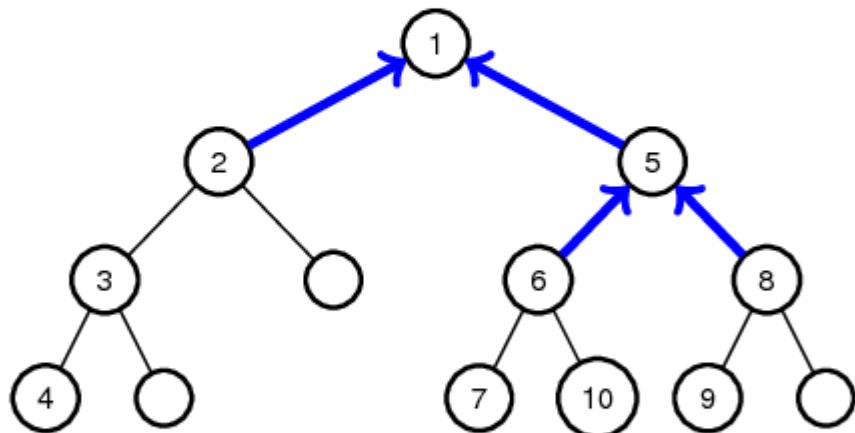
- Improved anytime behavior
- Incrementality  
without memory overhead



# HBFS

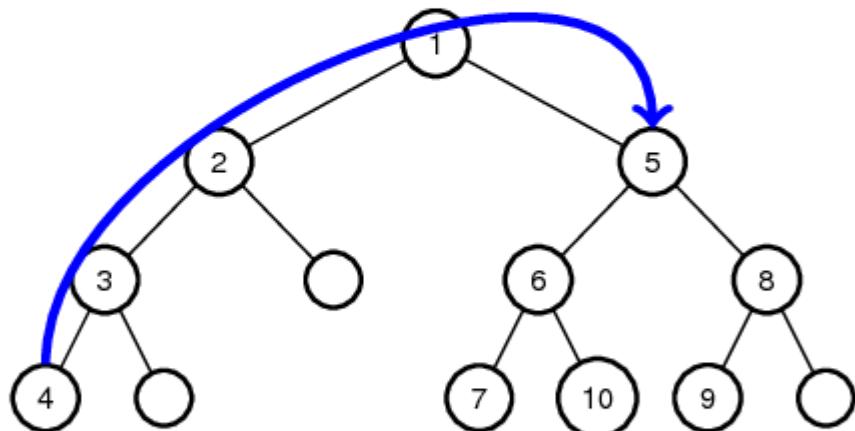
BFS with DFS probes\*

- Improved anytime behavior
- Incrementality without memory overhead
- Lower bounds



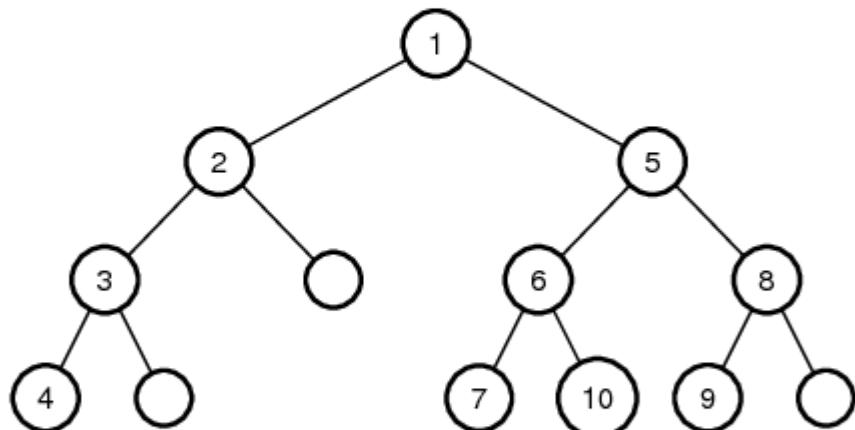
# HBFS

BFS with DFS probes\*



- Improved anytime behavior
- Incrementality without memory overhead
- Lower bounds
- Some of the advantages of restarting

# HBFS



BFS with DFS probes\*

- Improved anytime behavior
- Incrementality without memory overhead
- Lower bounds
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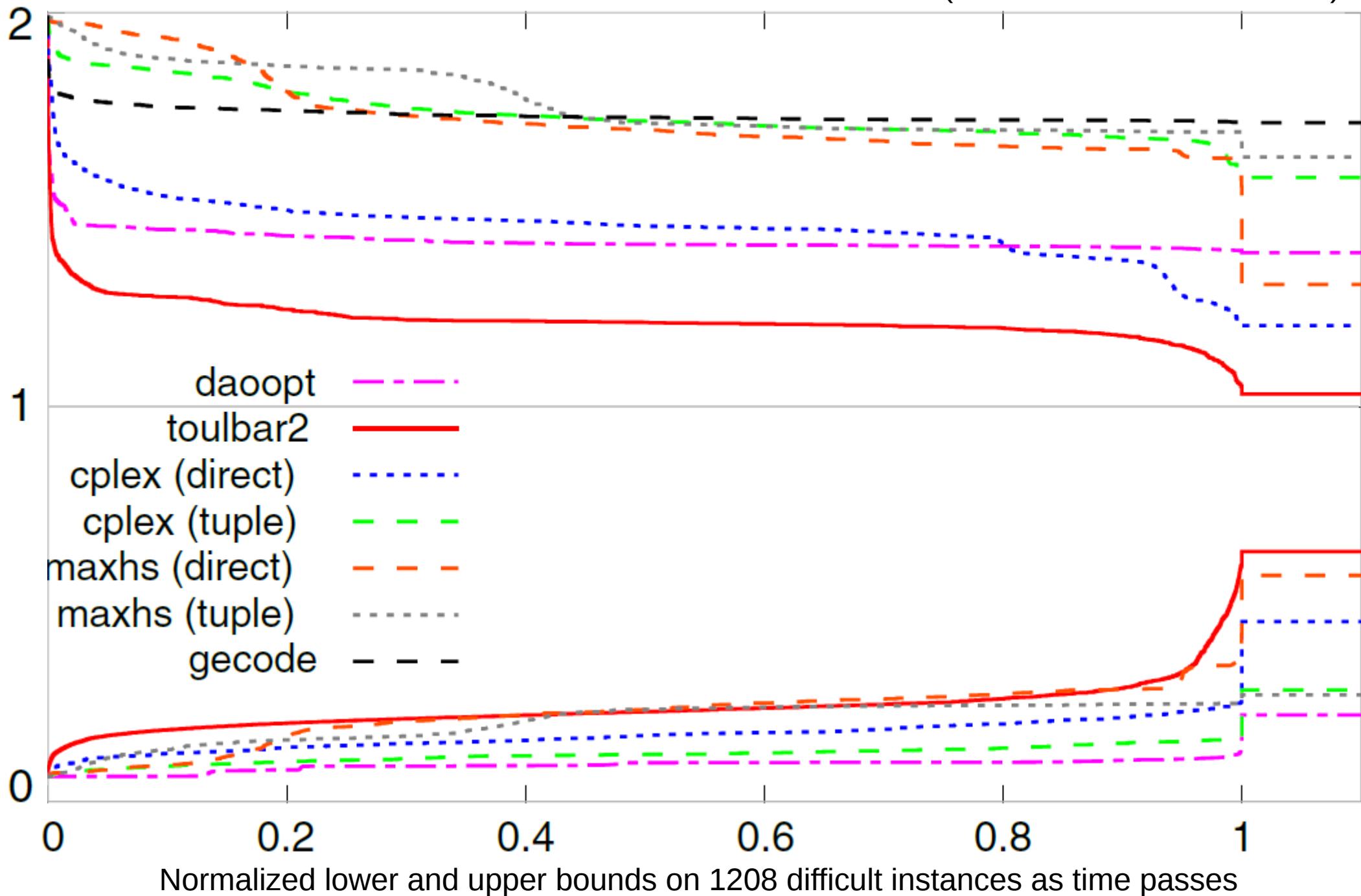
\* With adaptive heuristic for probe size

# Benchmark

- MRF: Probabilistic Inference Challenge 2011 (uai format)
- CVPR: Computer Vision and Pattern Recognition OpenGM2 (uai)
- CFN: MaxCSP 2008 Competition and CFLib (wcsp format)
- WPMS: Weighted Partial MaxSAT Evaluation 2013 (wcnf format)
- CP: MiniZinc Challenge 2012 & 2013 (minizinc format)

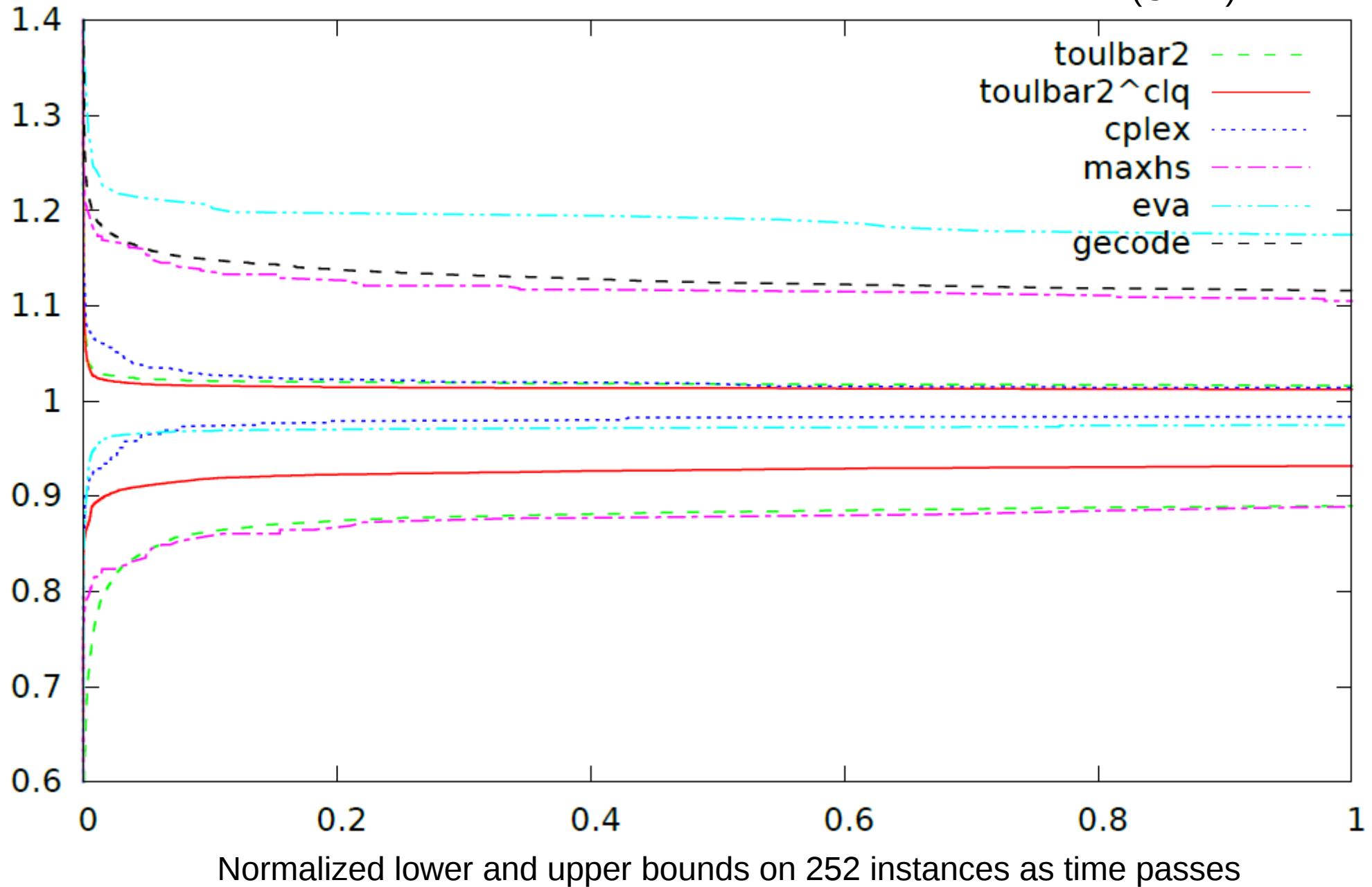
Number of instances and their total compressed (gzipped) size:

Benchmark	Nb.	UAI	WCSP	LP(direct)	LP(tuple)	WCNF(direct)	WCNF(tuple)	MINIZINC
MRF	319	187MB	475MB	2.4G	2.0GB	518MB	2.9GB	473MB
CVPR	1461	430MB	557MB	9.8GB	11GB	3.0GB	15GB	N/A
CFN	281	43MB	122MB	300MB	3.5GB	389MB	5.7GB	69MB
MaxCSP	503	13MB	24MB	311MB	660MB	73MB	999MB	29MB
WPMS	427	N/A	387MB	433MB	N/A	717MB	N/A	631MB
CP	35	7.5MB	597MB	499MB	1.2GB	378MB	1.9GB	21KB
Total	3026	0.68G	2.2G	14G	18G	5G	27G	1.2G

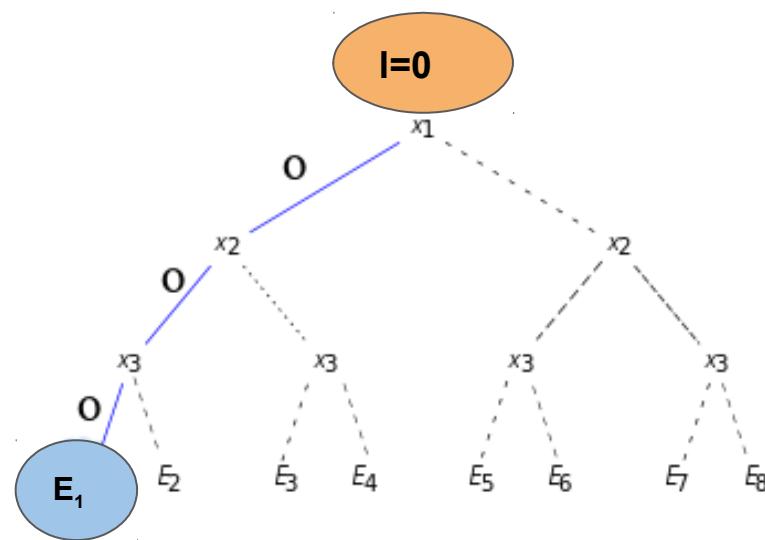


# Results exploiting cliques

(CP17)



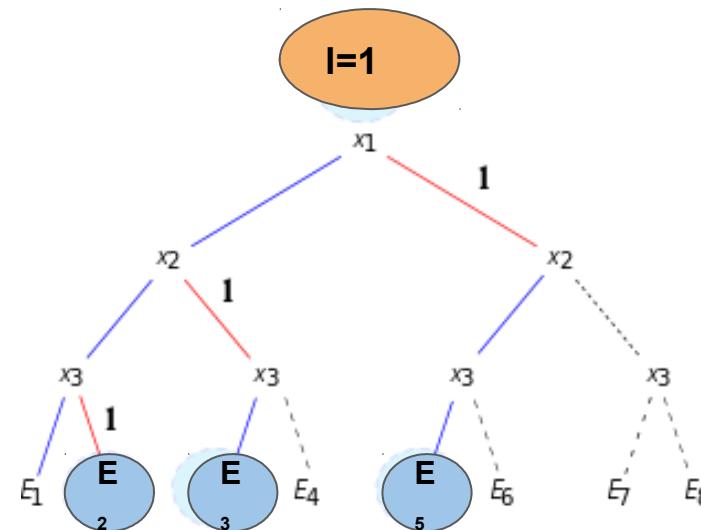
# Limited Discrepancy Search (*Ginsberg 95*)



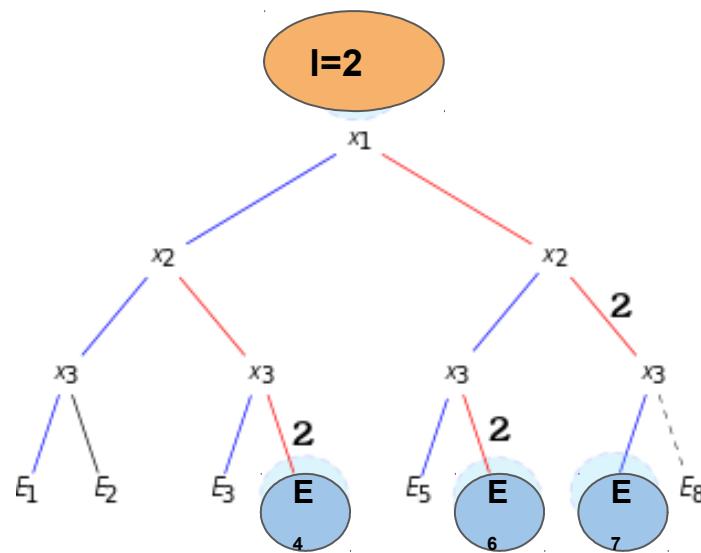
- Small example with 3 variables and 2 values per domain

# Limited Discrepancy Search

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# Limited Discrepancy Search (*Ginsberg 95*)

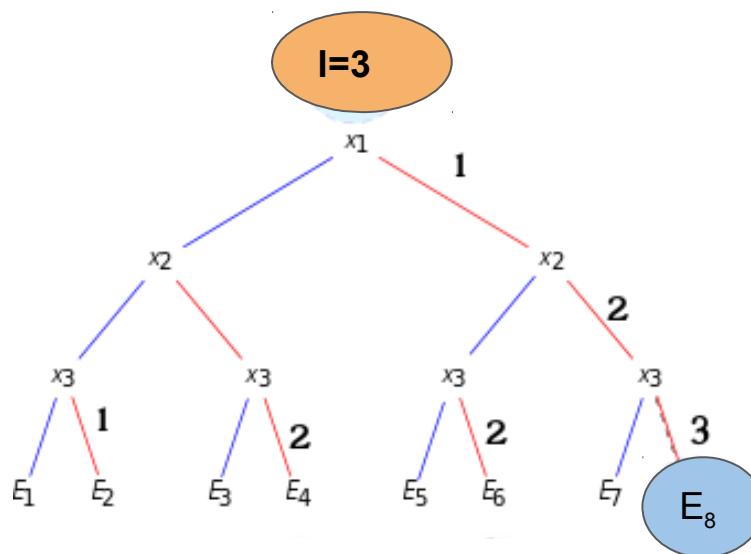
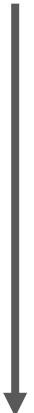


# Limited Discrepancy Search (*Ginsberg 95*)

$$l_{\max} = n * (d - 1) :$$

in this case,  $l_{\max} = 3 * (2 - 1) = 3$

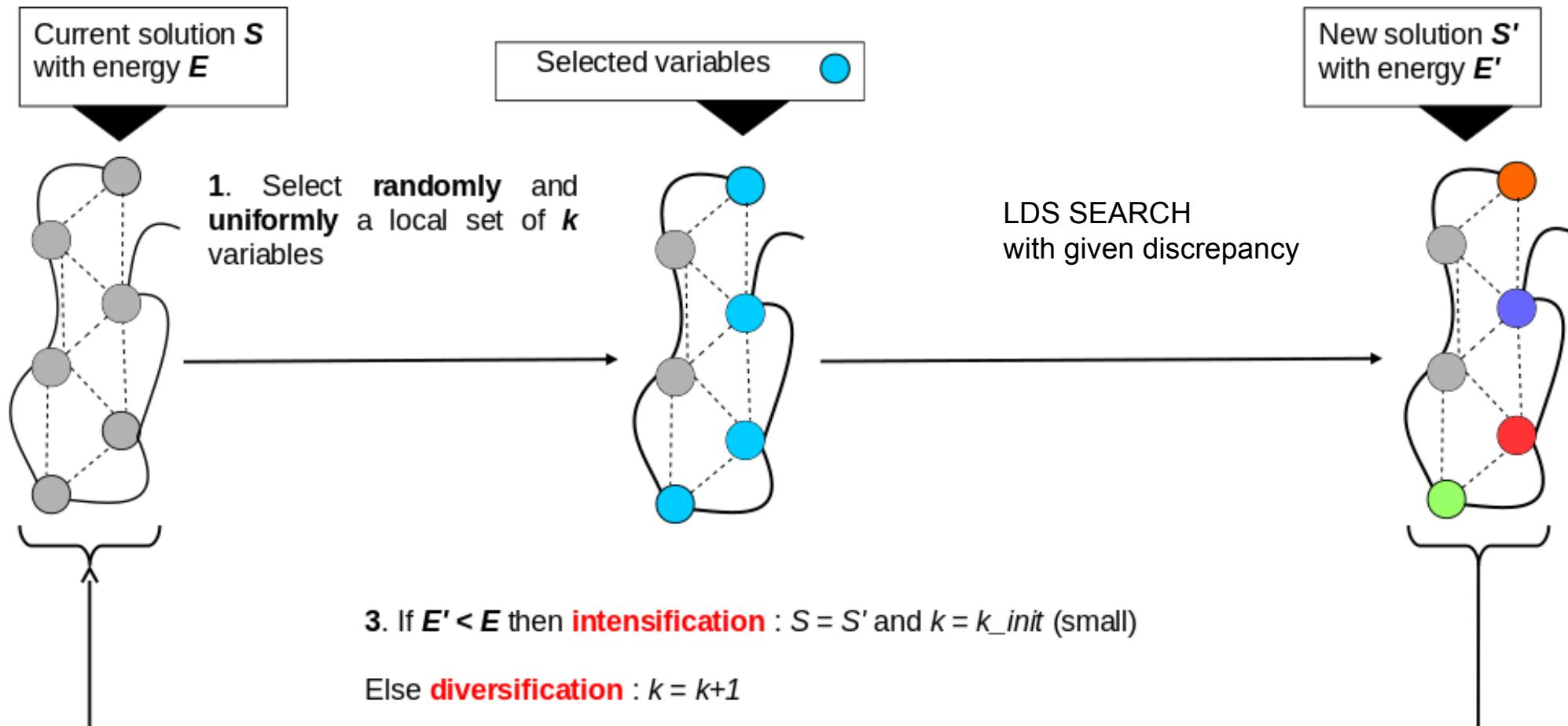
Full exploration



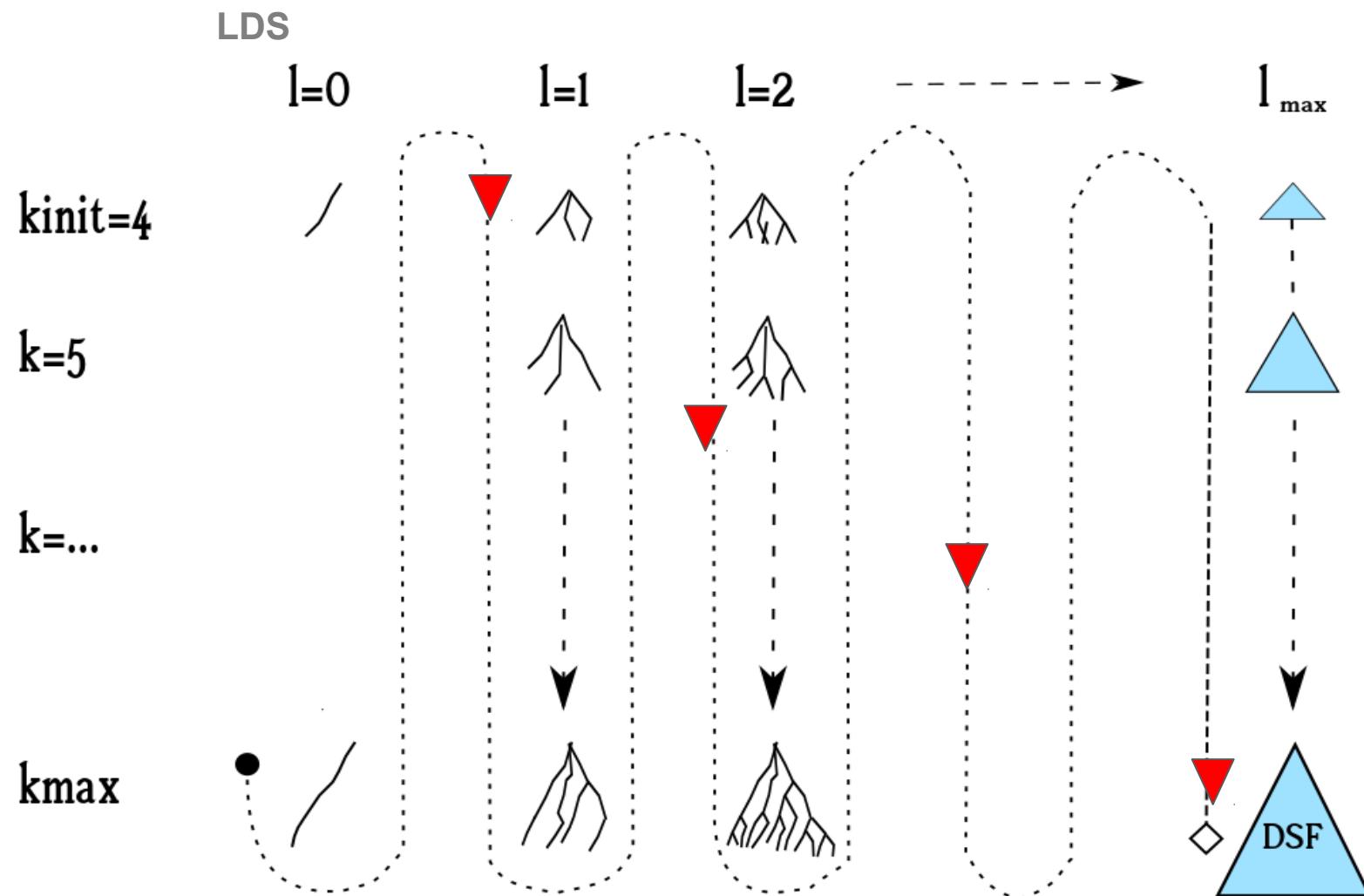
$l=3 \Rightarrow$  optimality proof

In practice, it occurs before  $l_{\max}$  thanks to bounding and pruning

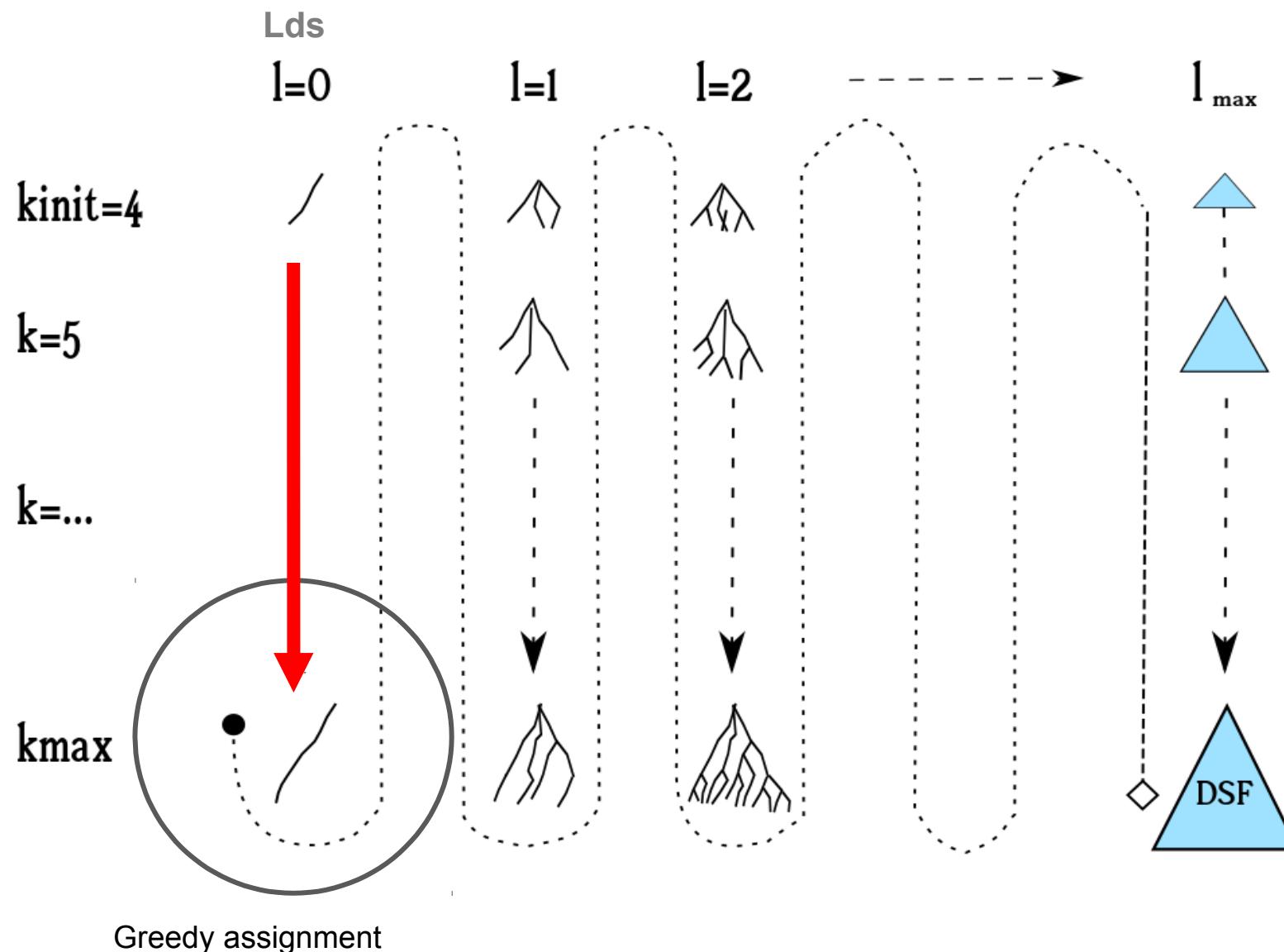
# Variable Neighborhood Search (Hansen 97)



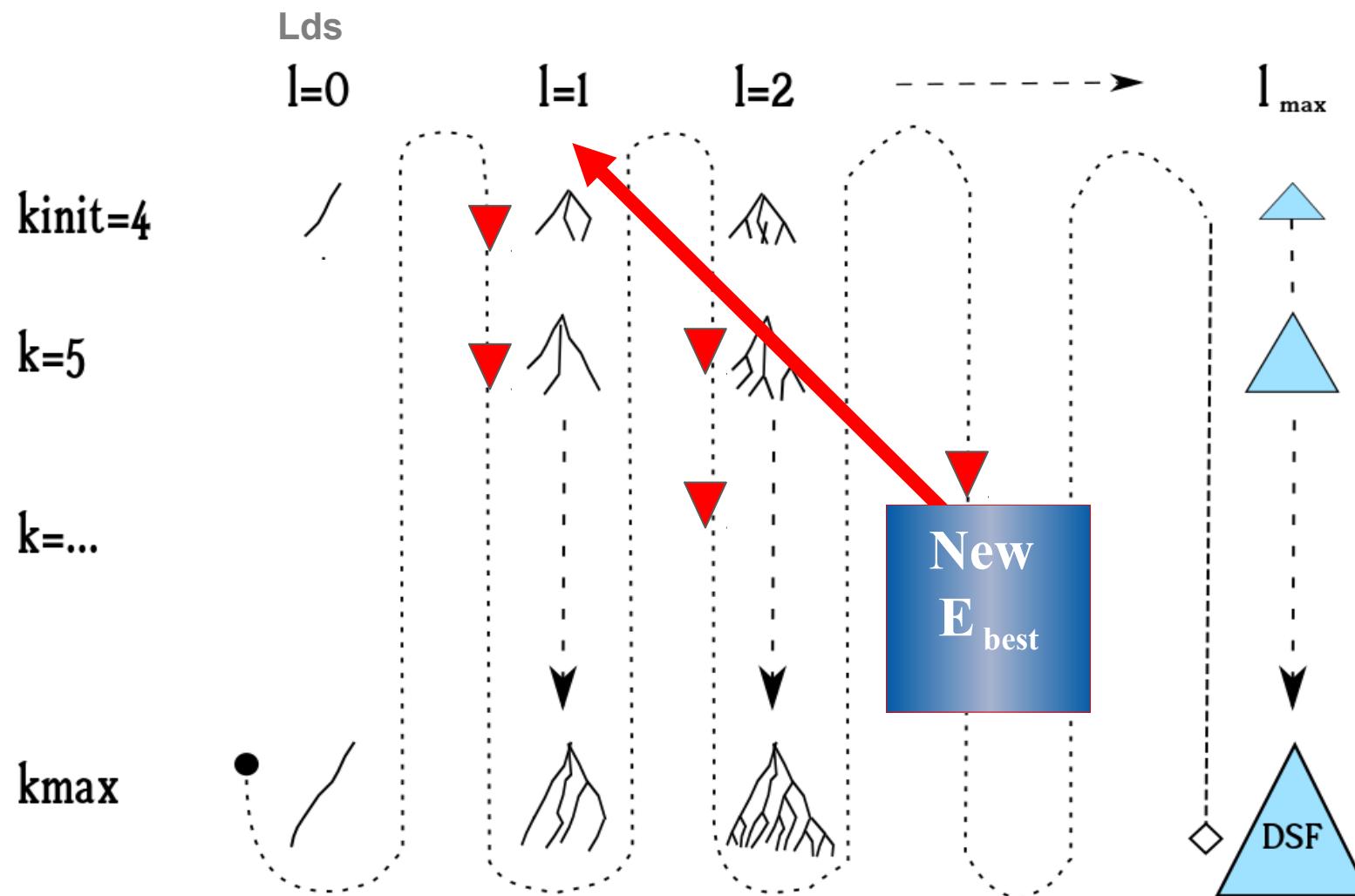
# UDGVNS : Exploration of both k and l dimensions



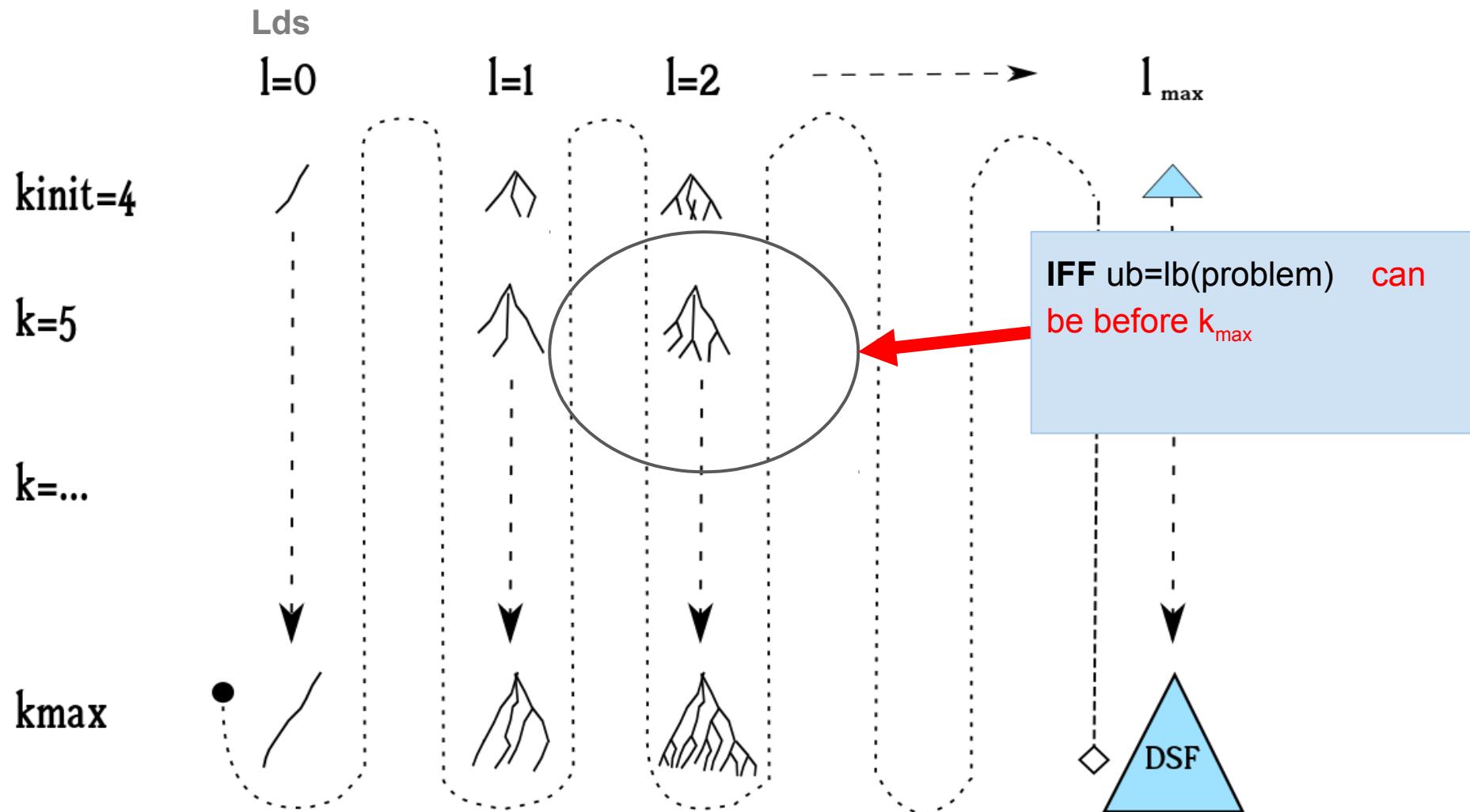
# Step 1 : Initial solution



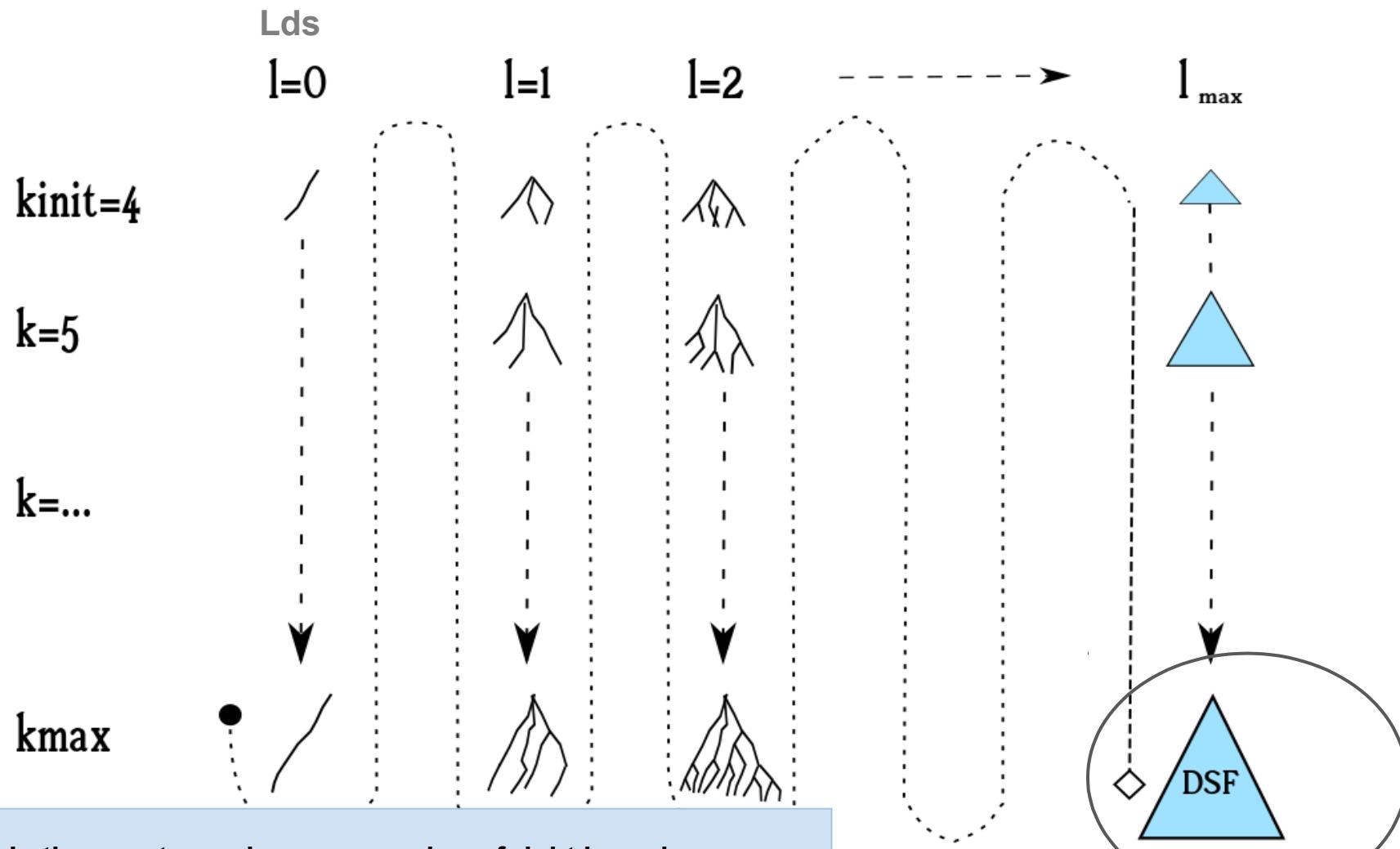
# NEW SOLUTION WITH BETTER E → RESTART



# Proof of Optimality

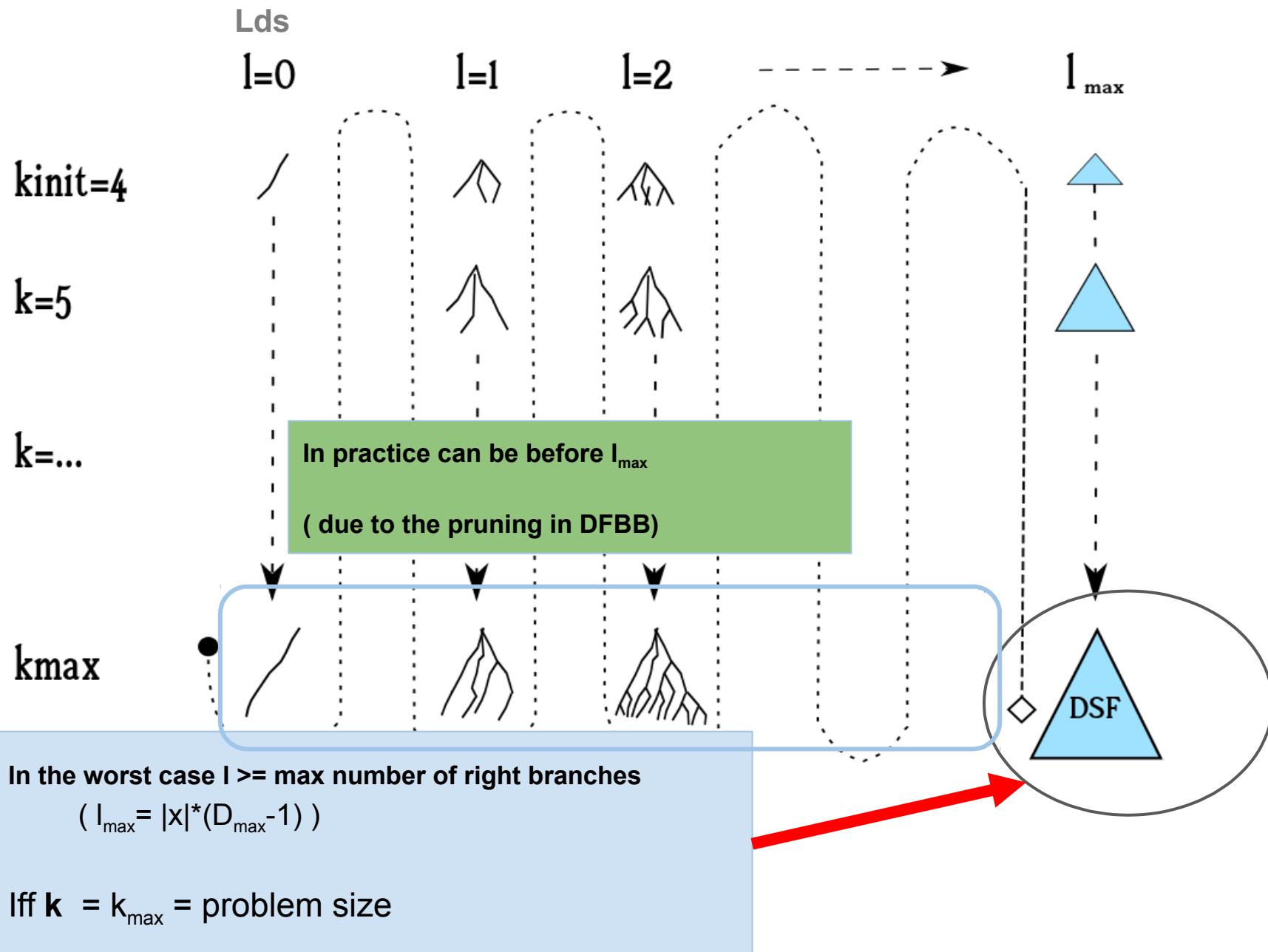


# Proof of Optimality



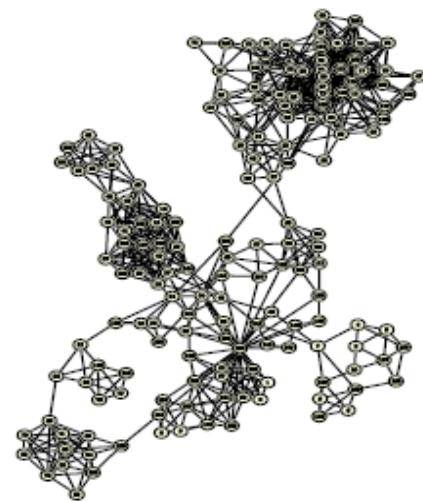
Iff  $k = k_{max} = \text{problem size}$

# Proof of Optimality

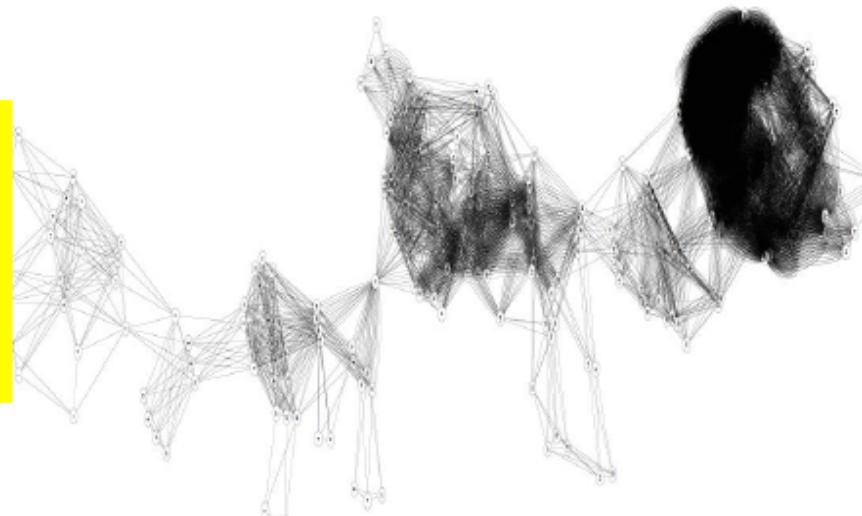


# Neighborhoods using problem structure

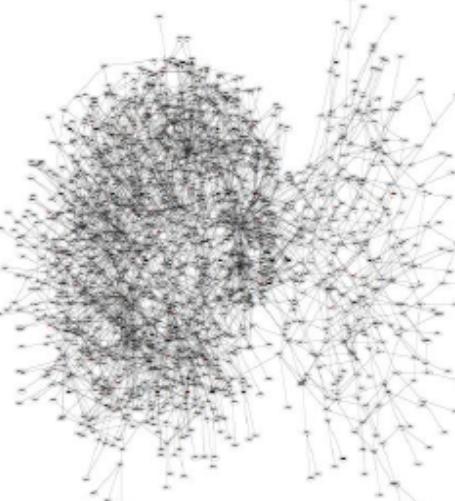
Radio  
Link  
Frequency  
Assignment



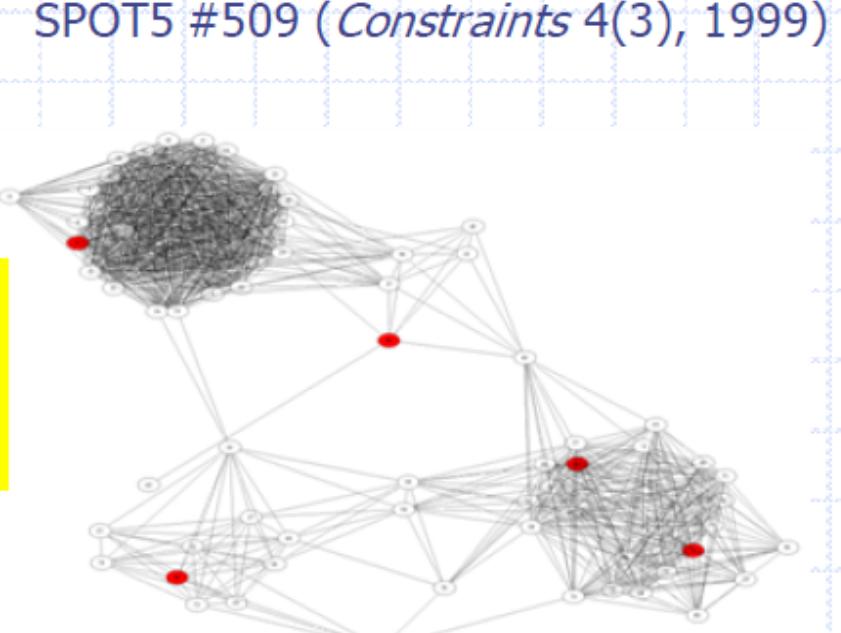
Earth  
Observation  
Satellite  
Management



Mendelian  
Error  
Detection



Tag  
SNP  
Selection

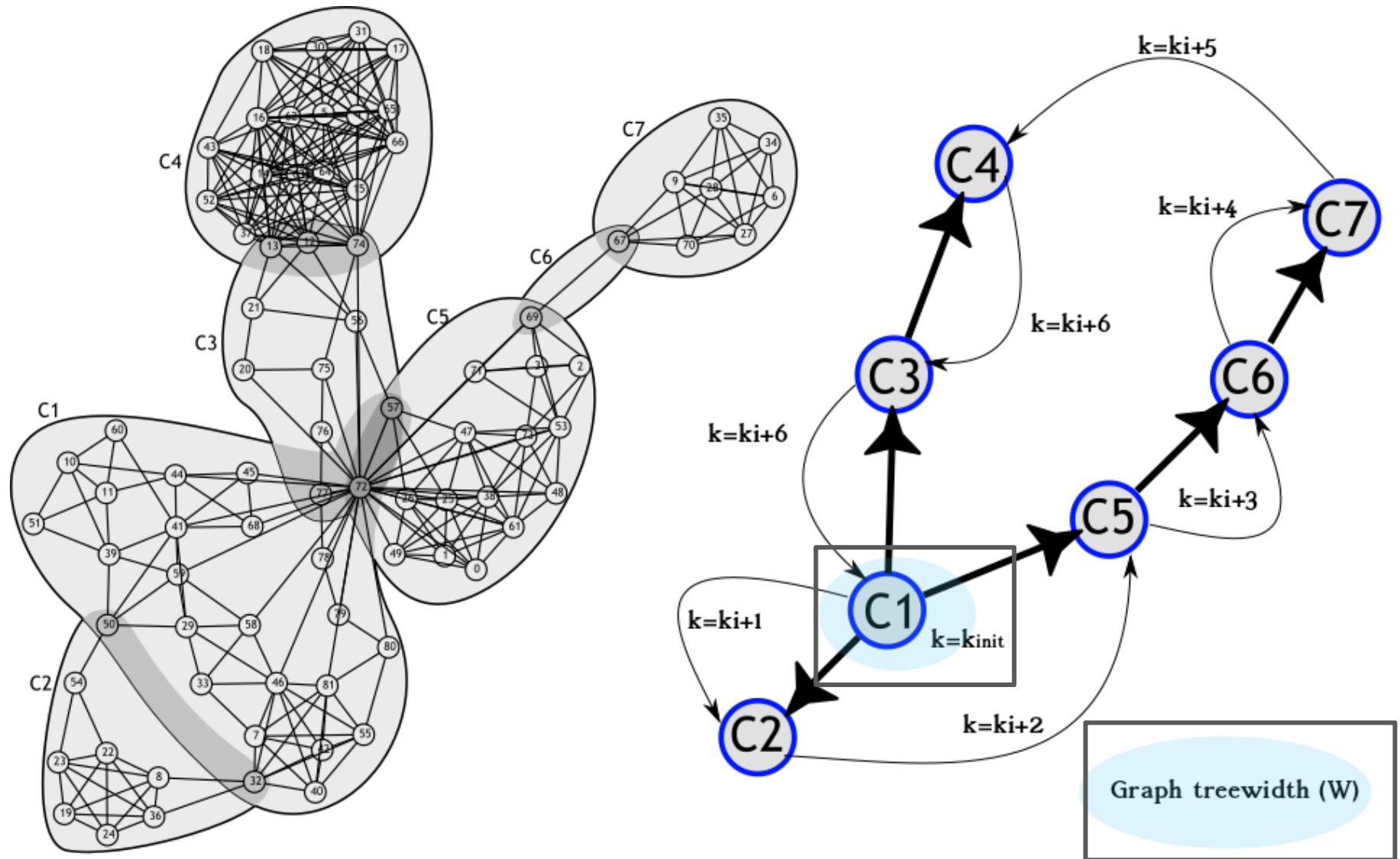


CELAR SCEN-07r  
(*Constraints* 4(1), 1999)

langladeM7 sheep pedigree  
(*Constraints* 13(1), 2008)

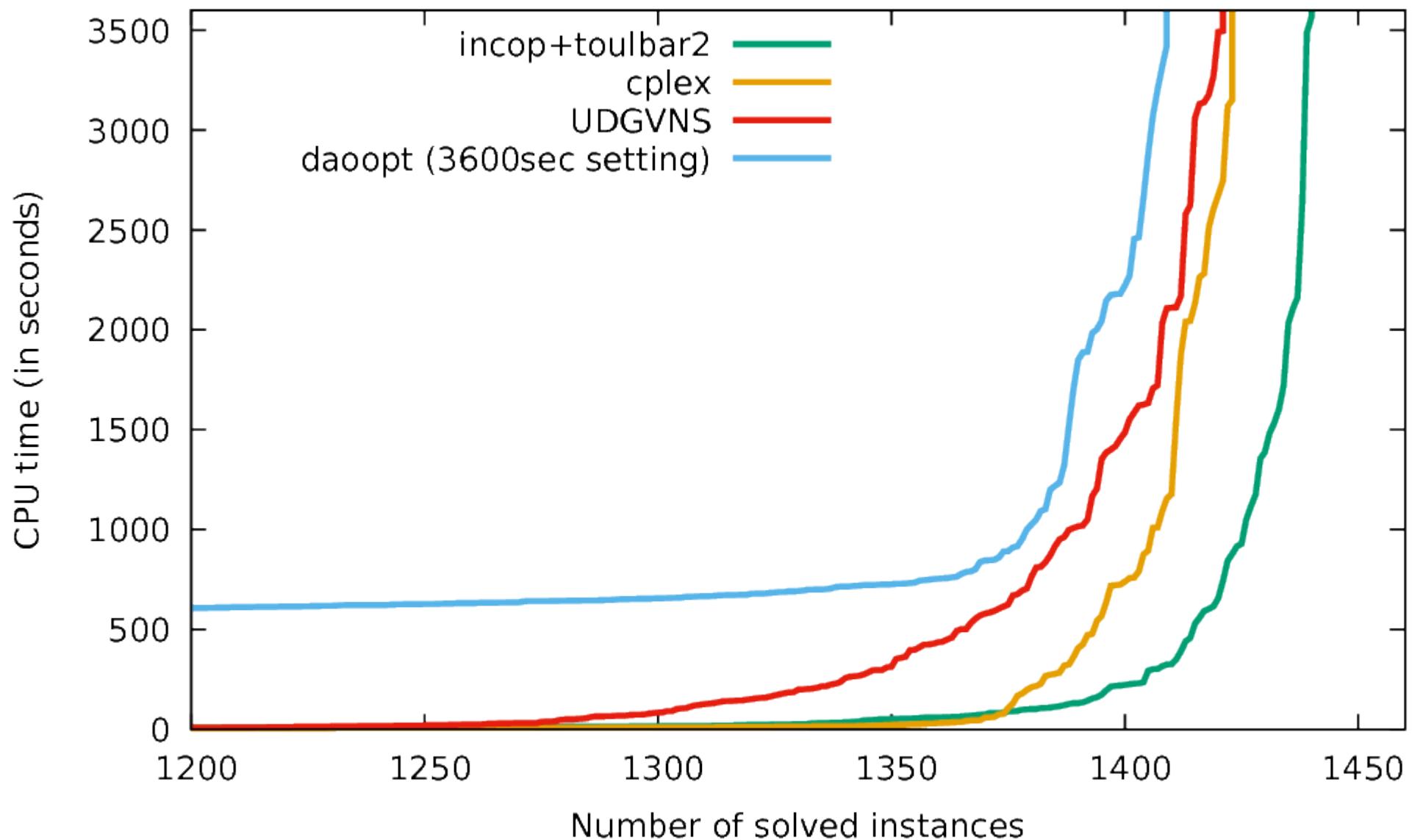
HapMap chr01  $r^2 \geq 0.8$  #14481  
(*Bioinformatics* 22(2), 2006)

# Cluster visit in a topological order :



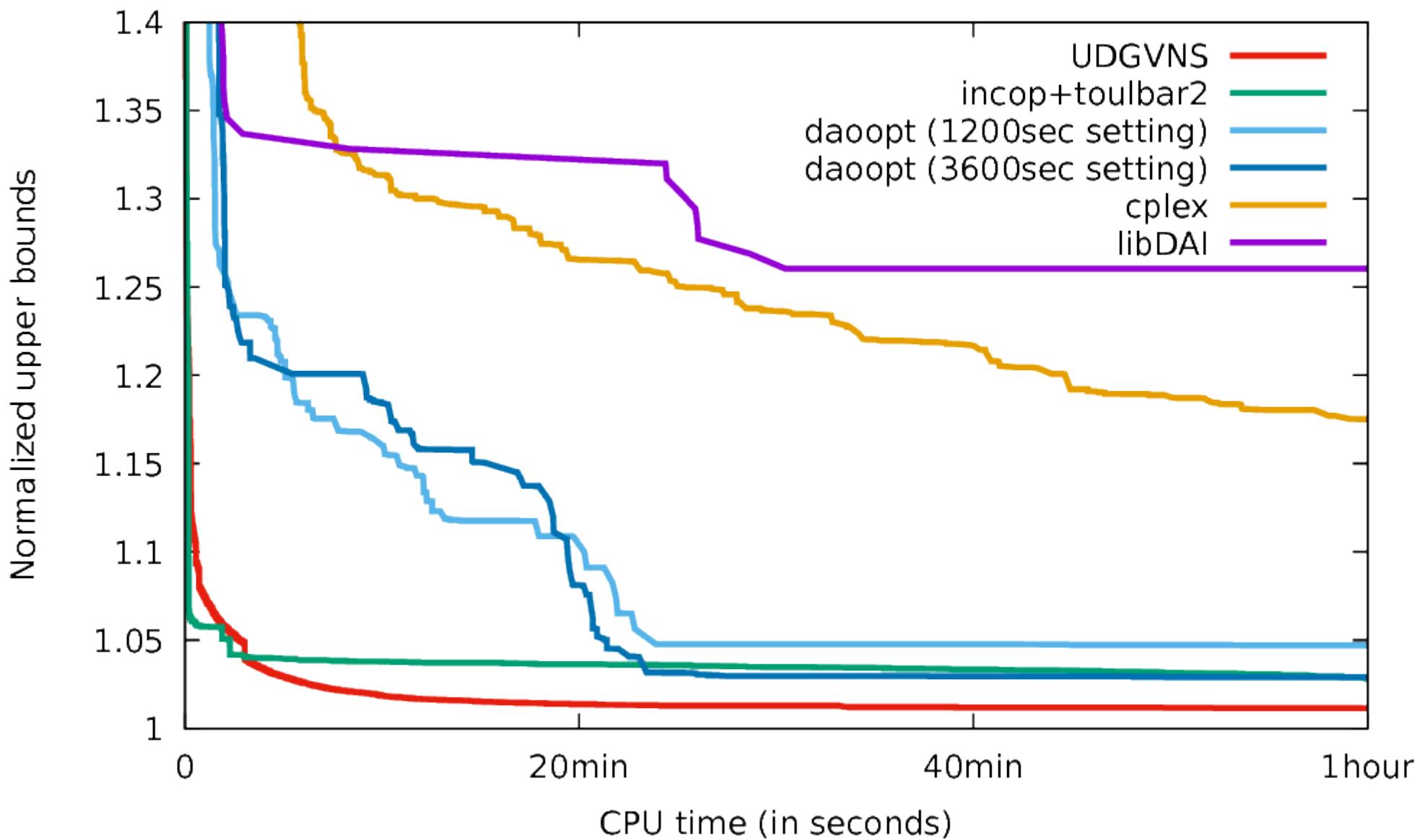
# Results

(UAI17)



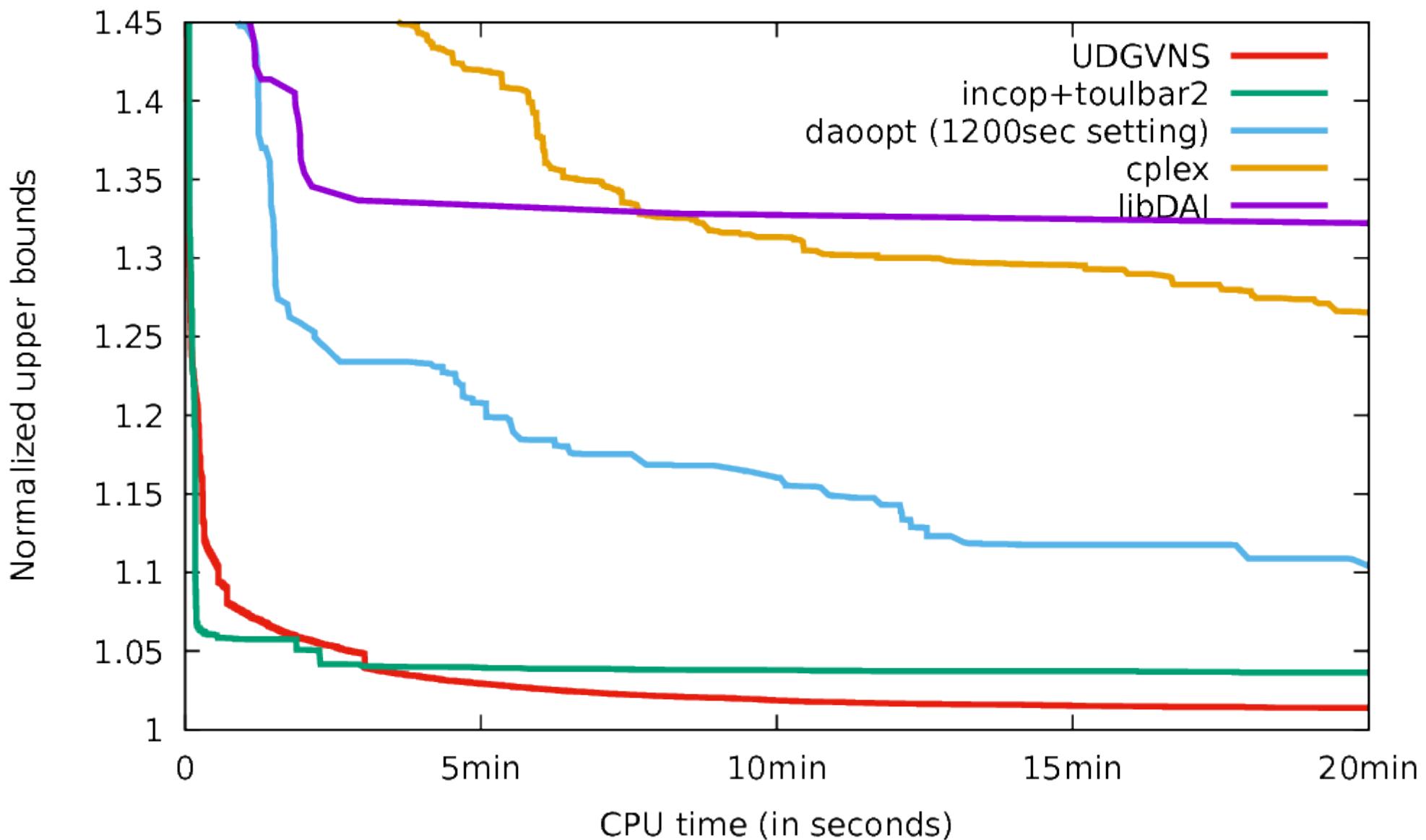
# Results

(UAI17)

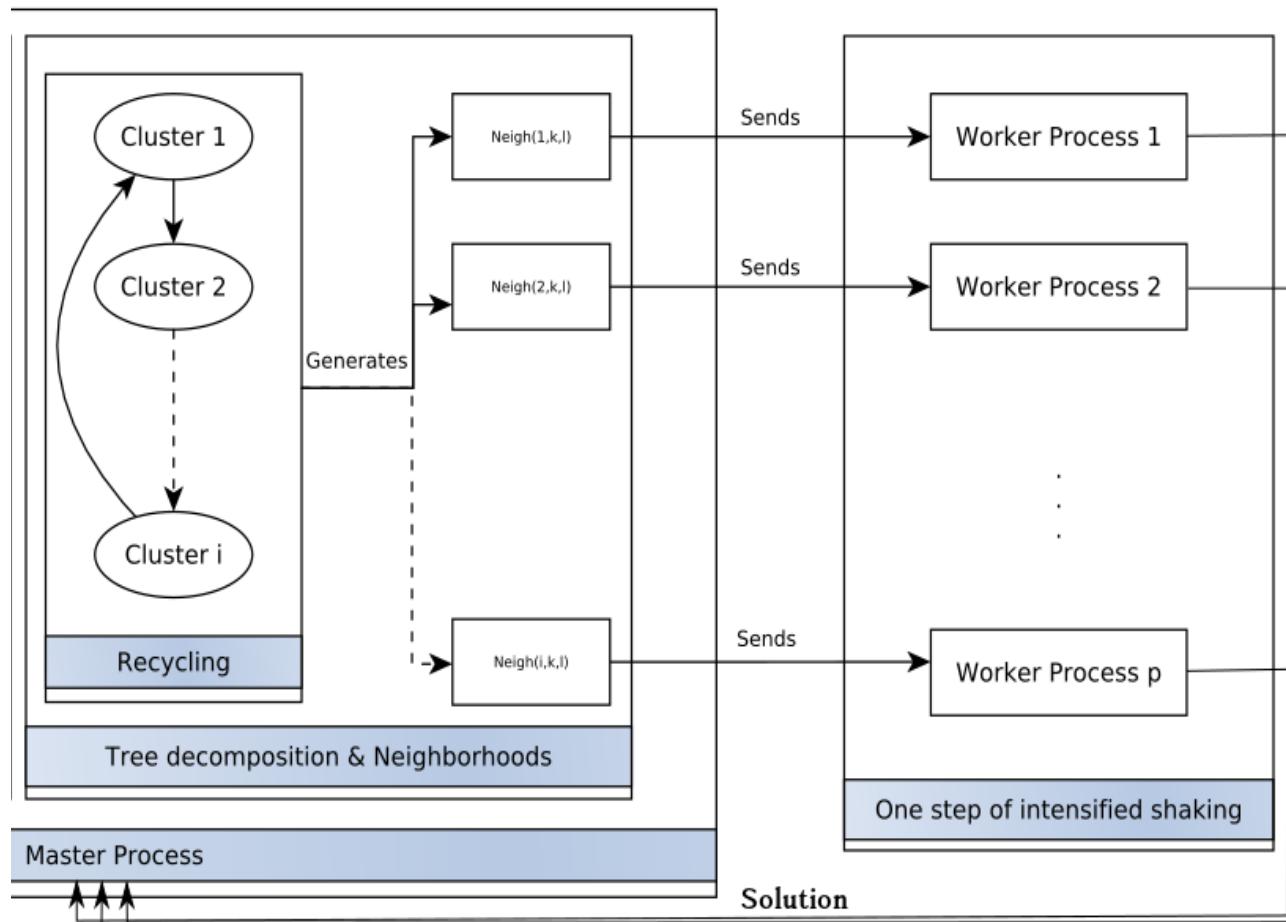


# Results

(UAI17)



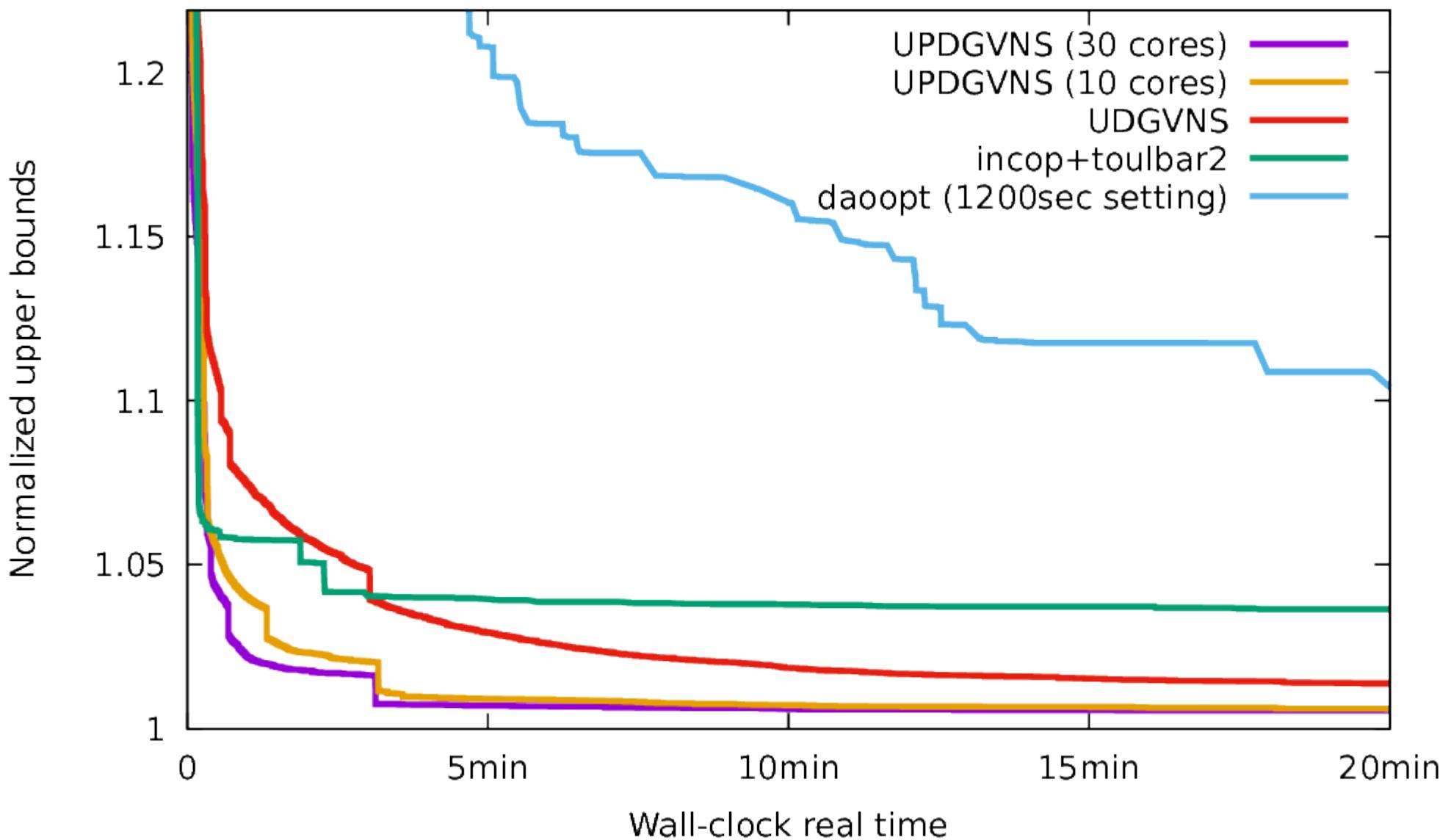
# Parallel VNS



*Unified Parallel Decomposition Guided VNS (UPDGVNS)*

# Results

(UAI17)



# References

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<https://arxiv.org/pdf/1506.08544.pdf>
- Hurley et al. *Multi-Language Evaluation of Exact Solvers in Graphical Model Discrete Optimization*. Constraints, 2016.  
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<https://github.com/toulbar2/toulbar2>

