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How Computers Break (Serious) Puzzles with logic and (a different breed of) learning

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How Computers Break (Serious) Puzzles

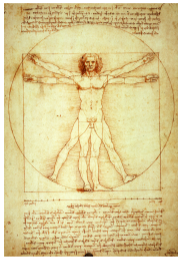
with logic and (a different breed of) learning

Thomas Schiex



February 2018

Académie des sciences, Paris, France

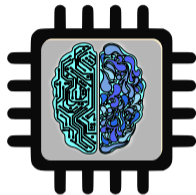


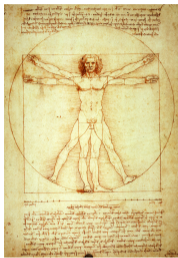
Human beings

- Easily rely on quick “intuitions” (ill-defined problems)
- Extreme rigor is painful and slow (logic/arithmetic)

Als (computers)

- Accessible to some “intuition” (problems defined by data)
- Fast and extreme rigor is the default (1 billion op./sec)



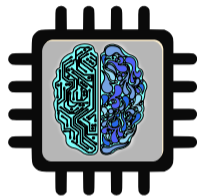


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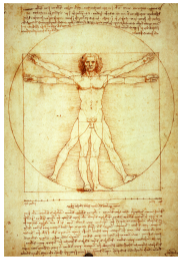
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It was expected that machines would show superhuman “logical reasoning” performances

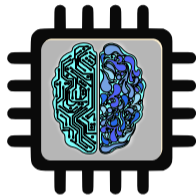


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1955: Newell & Simon “Logic Theorist” proved 38 of the 52 theorems in the *Principia Mathematica* (Russel and Whitehead), and even corrected a proof in it.

NP-hard problems

(Cook-Levin, 1970s)

- Some problems seems intrisically hard (for AIs at least)
- *Worst case asymptotic* exponential time ($P \neq NP$)

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$n^2 \times n^2$ Sudoku

- NP-complete, 9×9 : 10^{80} cases
- 10^{51} ages of the universe to examine them all
- Fast brute force will fail

							1	
				2				3
			4					
						5		
4		1	6					
		7	1					
	5					2		
				8			4	
	3		9	1				

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	E	1	6				D	7	F	G	B	
B			7			A	1	9	D		4	
A		2	C		1	B					E	
	7				G	F	2				3	
B			1								7	
9	A			6		1	E			G	B	
	8	C				5	7				1	6
G				2	A	F		3	5			C
E	5	G				B			9			
C			4			8	5	6				
	9	F	3	1	C		D		E			
8	4		5		F	D	3	G	1			6
	C			7	6				3	F		G
5		D	F			C		8		9		
F	2					3		4				
	6							D		C		5

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B	A	P	L		C	M		1	J	H	4	7									
		G	F	A	O	7	9	3	L	C	P	M	B	K							
N	M		I			4	P		6	A	E	F	G	C	D	3					
5	K							4	D	L	O	6									
C	H	1		D	F	G	N		B	E						9	P				
7	A							6	4				G	H		2	1	F			
G		J	N	L	E	B			B	C	7	1	P	F		H	O				
2			6	K			J	1	P	G	N		A	B	3	M					
	F	I	B	6	C	7	3	D	O	N	H		M	K	E	P					
	3		P					1	2	F	K	J				9	N				
L			K	C	9	N	E	6	H	A	B	M	3	B			2	7	J		
N	3	B	M						D		P		K						O		
	S	B	I	O		K	9			3	A		J	2	6	L					
P			F		L		6				4	G	H	K							
1	B	G			2	5	4	L	J	K	I	6	C	B	D	E			9		
		F	6			4	7	9	H	G			8		C						
	C	L	2	1			N	F	K	O	7	D	G	3	I	9	E				
	P	A	4	H		L	E	1	J	2				C	O				B		
	9	D		8	C	E	L	M	P			F	I	B	1	K			7		
		B	B	M		D	6			7	C								G	H	
3	G									H	6	O	M	C	E				I	P	1
				1	L	P	3		B							K				F	
	C	B	2		5	O	J	G	A		F	9				3				4	B
D	6			M	3		G	8	2	5	9	H	1		A	O					
	H	I	O	7	9			J		E				5		L	6			M	

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- Fast brute force will fail
- Can be solved in milliseconds

						1	
				2			3
			4				
					5		
4	1	6					
		7	1				
	5				2		
				8		4	
3		9	1				

Logic

- We have a set of variables

From a well defined problem to a solution
(Sudoku cells contain a number from 1 to 9)

x_1

x_2

x_3

x_4

x_5

x_6

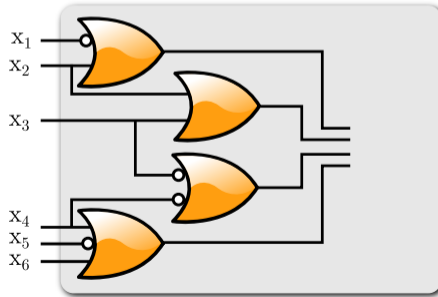
Logic

- We have a set of variables
- We have a set of properties on these variables

From a well defined problem to a solution

(Sudoku cells contain a number from 1 to 9)

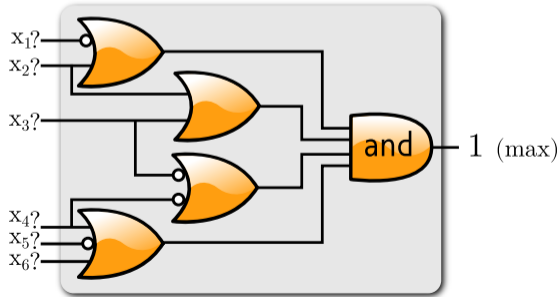
(all different rows, columns, super-cells)



Logic

- We have a set of variables (Sudoku cells contain a number from 1 to 9)
- We have a set of properties on these variables (all different rows, columns, super-cells)
- We want to find an input that satisfies all properties (or prove none exists: refutation).

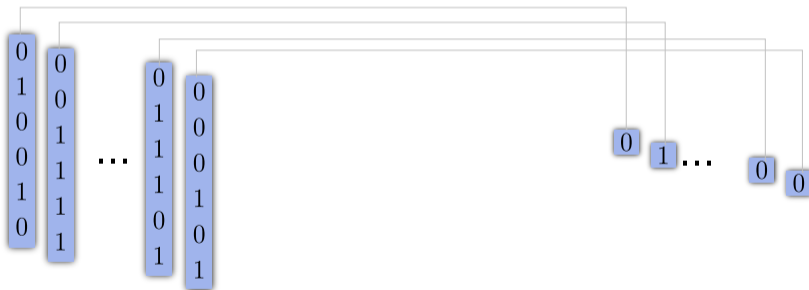
From a well defined problem to a solution



Intuition (DL)

From examples to a classifier

- We have a set of digital inputs (in \mathbb{B}^n) and output (class: one bit).
- We want a function that best predicts seen (and unseen) data in most cases.



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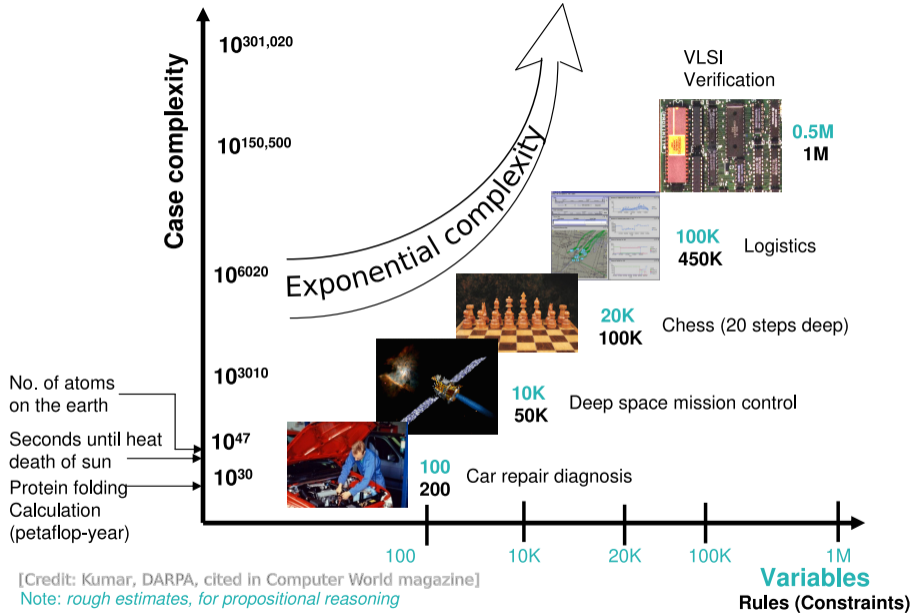


Technological progress

- Increasingly complex useful objects planes, computers, software, cars, Als
- That must be highly reliable (lives at stake)
- We cannot fully get them under control anymore

Increasing system complexity

- Hardware: Pentium FDIV bug (1994, 3.1 million transistors)
- Software: the Therac-25 (radiation-therapy) kills 6 patients
- Tesla cars: said to carry 100 millions lines of codes
- Convolutional NN: may have billions of parameters



SAT

- 1 A set of Boolean variables x_i
- 2 A set of clauses (*disjunction* of variables or negation of) $(\neg x_1 \vee x_7)$
- 3 Must satisfy all clauses (or prove impossible)
- 4 Semantics: defines a function from \mathbb{B}^n to \mathbb{B}

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Sudoku

- 1 cell (i, j) contains k x_{ijk} true
- 2 At least one number per cell i, j $(x_{ij1} \vee \dots \vee x_{ij9})$
- 3 At most one number per cell i, j $(\forall k > k' \neg x_{ijk} \vee \neg x_{ijk'})$
- 4 Cell (i, j) and (i, j') must be different $(\neg x_{ijk} \vee \neg x_{ij'k})$

More sophisticated/practical function description

- propositions over theories
- non Boolean variables
- numerical output

SAT Modulo Theory⁹

Constraint Satisfaction, Constraint Programming³⁰

Weighted MaxSAT²⁵/CSP,⁵ Graphical models¹⁸

SAT is the simplest

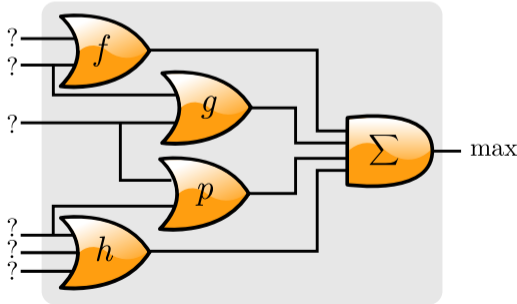
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NP-complete: can express all NP-complete problems

- the logical puzzles you like (Sudoku, Nonograms...)
- or not (configuration, scheduling, test pattern generation...)
- robot planning
- digital circuit verification (Bounded Model Checking)
- or software verification (FOL, grounding, abstraction)











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NP-complete, so intractable

Standard argument for less realistic problem reformulation, heuristics or stochastic search

NP-complete: can express all NP-complete problems

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- or not (configuration, scheduling, test pattern generation...)   
- robot planning (Rosetta-Philae probe plan, CP, LAAS/Toulouse) 
- digital circuit verification (Bounded Model Checking)  
- or software verification (FOL, grounding, abstraction)   

NP-complete, so intractable

Standard argument for less realistic problem reformulation, heuristics or stochastic search

Real SAT instances with millions of variables/clauses can be solved (with a proof)

```
p cnf 51639 368352
-1 7 0
-1 6 0
-1 5 0
-1 -4 0
-1 3 0
-1 2 0
-1 -8 0
-9 15 0
-9 14 0
-9 13 0
-9 -12 0
-9 11 0
-9 10 0
-9 -16 0
```

51,639 variables, 368,352
constraints

$\neg x_1 \vee x_7$

$\neg x_1 \vee x_6$

...

185 -9 0

185 -1 0

177 169 161 153 145 137 129

121 113 105 97 89 81 73 65 57

49 41 33 25 17 9 1 -185 0

186 -187 0

186 -188 0

...

$(x_{177} \vee x_{169} \vee x_{161} \vee x_{153} \vee \dots \vee$
 $x_{17} \vee x_9 \vee x_1 \vee \neg x_{185})$

10236 -10050 0
10236 -10051 0
10236 -10235 0
10008 10009 10010 10011 10012 10013 10014 10015 10016 10017 10018
10019 10020 10021 10022 10023 10024 10025 10026 10027 10028 10029
10030 10031 10032 10033 10034 10035 10036 10037 10086 10087 10088
10089 10090 10091 10092 10093 10094 10095 10096 10097 10098 10099
10100 10101 10102 10103 10104 10105 10106 10107 10108 -55 -54 53 -52
-51 50 10047 10048 10049 10050 10051 10235 -10236 0
10237 -10008 0
10237 -10009 0
10237 -10010 0
...

```
-7 260 0
7 -260 0
1072 1070 0
-15 -14 -13 -12 -11 -10 0
-15 -14 -13 -12 -11 10 0
-15 -14 -13 -12 11 -10 0
-15 -14 -13 -12 11 10 0
-7 -6 -5 -4 -3 -2 0
-7 -6 -5 -4 -3 2 0
-7 -6 -5 -4 3 -2 0
-7 -6 -5 -4 3 2 0
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```

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```

Search space

$$2^{50,000} \approx 3.1 \cdot 10^{15,051}$$

-7 260 0
7 -260 0
1072 1070 0
-15 -14 -13 -12 -11 -10 0
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-7 -6 -5 -4 3 2 0
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Solved in one second

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1072 1070 0
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-15 -14 -13 -12 11 -10 0
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-7 -6 -5 -4 -3 -2 0
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-7 -6 -5 -4 3 -2 0
-7 -6 -5 -4 3 2 0
185 0
```

Search space

$$2^{50,000} \approx 3.1 \cdot 10^{15,051}$$

Solved in one second

How does it work?

Consensus/Resolution (1960's)^{8,29}

$$\frac{(x_1 \vee \overbrace{l_1 \vee \dots \vee l_n}^L) \quad (\neg x_1 \vee \overbrace{r_1 \vee \dots \vee r_m}^R)}{(L \vee R)}$$

Boolean Constraint Propagation (BCP): unit clauses⁷

(L or R empty)

$$\frac{(x_1) \quad (\neg x_1 \vee \overbrace{r_1 \vee \cdots \vee r_m}^R)}{(R)}$$

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- This may create new unit clauses (propagation)

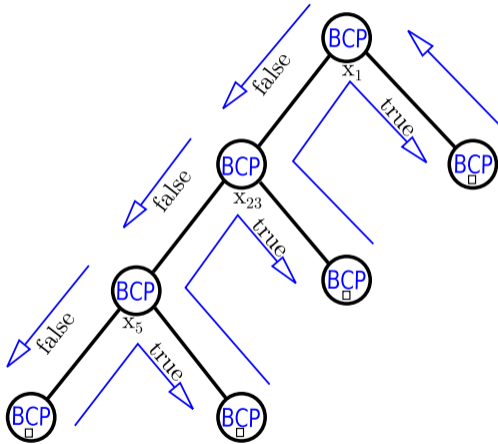
Boolean Constraint Propagation (BCP): unit clauses⁷

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- A clause is shortened by one literal
- This may create new unit clauses (propagation)
- If the empty clause \square appears: no solution

Logic: Try to guess and reconsider (DPLL⁷)



SAT state-of-the-art in 1990
Hundreds of variables
Thousands of clauses

Long line of research in “symbolic” Artificial Intelligence^{3,10,23,24,32}

- Trace back failure to guesses through propagation^a
- Do backward resolution from conflict
- Add a new implied clause to the set of clauses

^aRichard M Stallman and Gerald J Sussman. “Forward reasoning and dependency-directed backtracking in a system for computer-aided circuit analysis”. In: *Artificial intelligence* 9.2 (1977), pp. 135–196.

- Forces to reconsider an earlier guess
- Prevents re-failing for a related reason (safe generalization)

Learns a more effective formulation of the problem as it solves it

Intuition: Choose a variable and try to guess its value

Learning by “Activity based heuristics”²⁶

- On-line estimation of how often a variable is involved in recent clauses/failures
- Try guessing this variable first

Learns weak spots in the problem as it is solved

(safe)

A lot of free data and free code...

- International competitions (> 50,000 benchmarks with many real problems)
- Open source solvers (autocatalytic)



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Strong French presence



- Award winning solvers (Glucose,² toulbar2¹⁵)
- Constraint programming solver/startup (Choco)
- Strong presence in international conferences (# of accepted papers in CP⁴)

A conjecture in combinatorics

∞

When one splits \mathbb{N} in 2, one part must contain a Pythagorean triple

$$(a^2 = b^2 + c^2)$$

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No known proof, puzzled mathematicians for decades (one offered a 100 \$ reward)

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SAT solver proof^{14,22}

200TB proof, compressed to 86GB (stronger proof system)^a

^aOliver Kullmann. “The Science of Brute Force”. In: *Communications of the ACM* (2017).

Whether it's maths or not...

Size matters!

- Not only there exists true unprovable statements (in powerful enough consistent sets of axioms¹²⁾)
- There may be true provable statements we will never be able to prove because of their extremely long proofs²⁰



Biology

- Many discrete object ($\{A, T/U, G, C\}$, amino acids, genes, alleles, enzymes...)
- Lots of experimental data

Is it bio-compatible?

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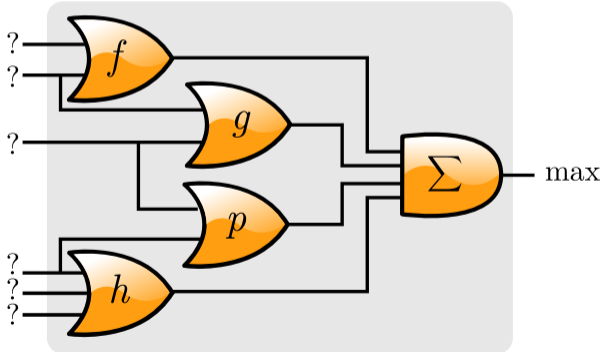
Exploiting Data + knowledge: Machine Learning

- (Stochastic) models can be built from knowledge and data
- And used to predict a “Most Likely/Optimal State” \Rightarrow easily NP-hard

Is it bio-compatible?

Biology

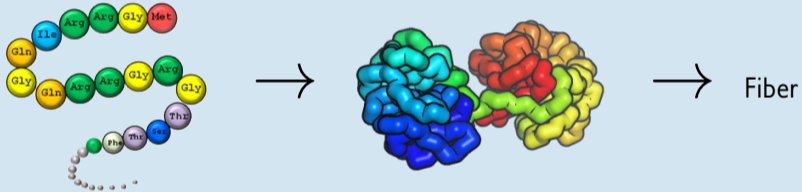
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Most active molecules of life

Sequence of “amino-acids”, each chosen among a set of 20 natural ones

Folding



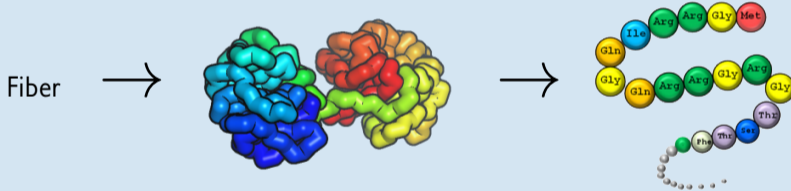
Transporter, binder, regulator, motor, catalyst...

Hemoglobine, TAL effector, ATPase, dehydrogenases...

Most active molecules of life

Sequence of “amino-acids”, each chosen among a set of 20 natural ones

Inverse folding



Transporter, binder, regulator, motor, catalyst...

Hemoglobine, TAL effector, ATPase, dehydrogenases...

Eco-friendly chemical/structural nano-agents

- New catalysts for biomass transformation (biofuels, food and feed, cosmetics...),
- New drugs for medicine
- New components for nanotechnologies

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20^n sequences!

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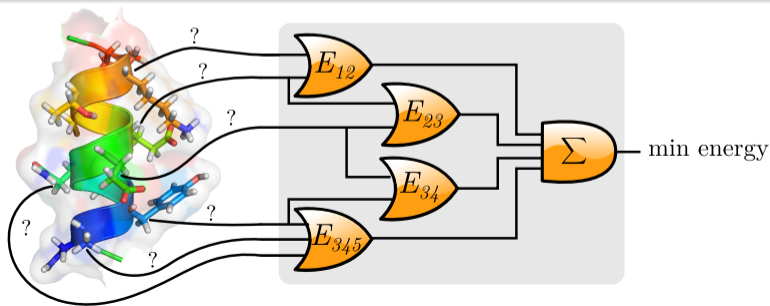
CPD: From bits to atoms

From information to functional matter

- mind blowing mass 3d printing-like capacities at atomic level (bacterias)
- structural and functional purposes (powerful origami)
- produced new folds,¹⁹ catalysts,³¹ nano-components³⁶

Ingredients

- Full atom model of a protein backbone (assumed to be rigid)
- Catalog of all 20 amino acids in different conformations (≈ 400 overall)
- Full atom energy function (bonds, electrostatics, solvent, statistics...)
- Maximum stability \equiv Minimum energy NP-hard²⁸



Large input (> 1GB)

NP-hard problem

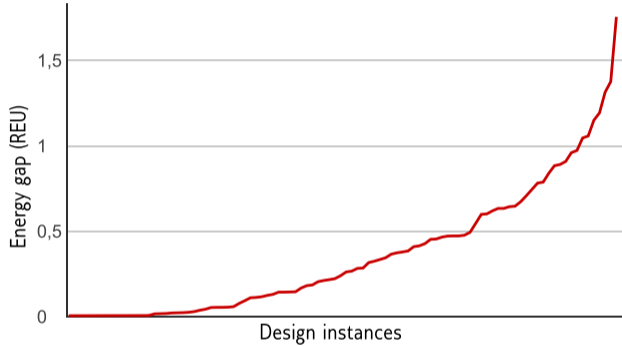
Toulbar2 is able to...

- provide a proven minimum energy solution
- exhaustively enumerate sequences close to it
- in spaces of size $> 10^{200}$



Showed that an highly tuned biased Monte Carlo increasingly fails to find the optimal sequence^a

^aDavid Simoncini et al. "Guaranteed Discrete Energy Optimization on Large Protein Design Problems". In: *Journal of Chemical Theory and Computation* 11.12 (2015), pp. 5980–5989. DOI: 10.1021/acs.jctc.5b00594.

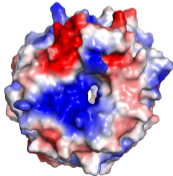
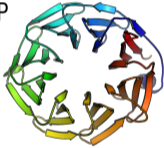


Asymptote: **Size matters!**


Asymptotic convergence can be arbitrarily slow...

C8 pseudo-symmetric 20VP symmetrized into a nano-component

20VP

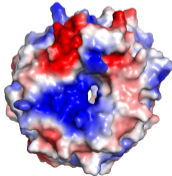
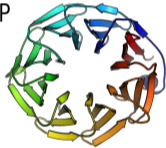


C8 pseudo-symmetric 20VP symmetrized into a nano-component

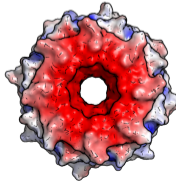
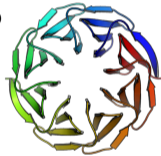
-  Tako: (R)evolution + Rosetta/talaris14

8 fold


20VP



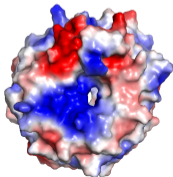
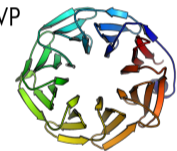
Tako



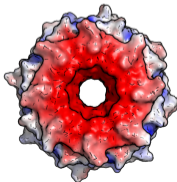
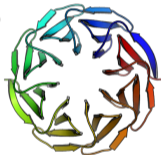
C8 pseudo-symmetric 20VP symmetrized into a nano-component

-  Tako: (R)evolution + Rosetta/talaris14 8 fold
-  Ika: toulbar2 + talaris14 4 fold

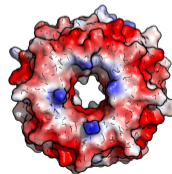
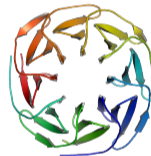
20VP



Tako



Ika



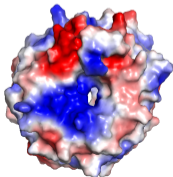
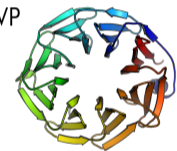
C8 pseudo-symmetric 20VP symmetrized into a nano-component



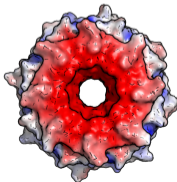
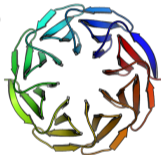
lka: toulbar2 + talaris14

4 fold

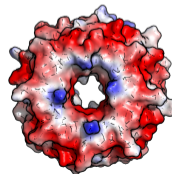
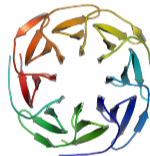
20VP

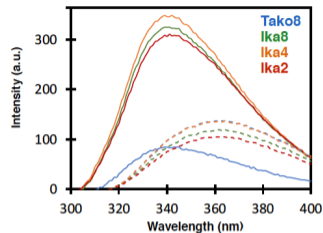
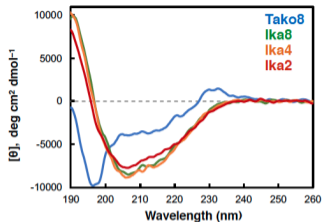
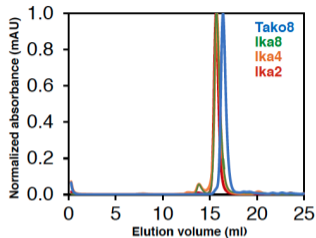


Tako



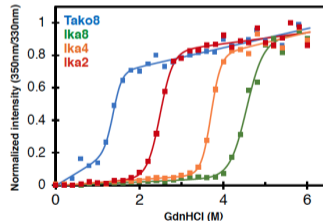
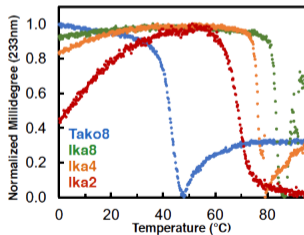
lka





Assemble as 8-bladed propeller

- Ika* more stable than Tako8
- Temperature
- Chemical denaturation



Asymptotes: size matters

NP is not exactly as we tend to think

- AIs have made drastic progress in their logical capacities
- This progress also comes from (gradient-free) learning
- More progress is needed to supplement our limited human capacities

Synergies between Logic and Intuition

- Logic can analyze and exploit learnt models (not only Neural Nets)
- Intuition can help logic without tainting it (guidance)

Al/toulbar2

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G. Katsirelos (INRA)
M. Zytnicki (PhD, INRA)
D. Allouche (INRA)
H. Nguyen (PhD, INRA)
M. Cooper (IRIT, Toulouse)
J. Larrosa (UPC, Spain)
F. Heras (UPC, Spain)
M. Sanchez (Spain)
E. Rollon (UPC, Spain)
P. Meseguer (CSIC, Spain)
G. Verfaillie (ONERA, ret.)
JH. Lee (CU. Hong Kong)
C. Bessiere (LIMM, Montpellier)
JP. Métivier (GREYC, Caen)
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Protein Design

A. Voet (KU Leuven)
D. Simoncini (INSA, Toulouse)
S. Barbe (INSA, Toulouse)
S. Traoré (PhD, CEA)
C. Viricel (PhD)
RosettaCommons (U. Washington)
W. Sheffler (U. Washington)
PyRosetta (U. John Hopkins)
B. Donald (U. North Carolina)
K. Roberts (U. North Carolina)
T. Simonson (Polytechnique)
J. Cortes (LAAS/CNRS)

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We do not understand the sources of their efficiency

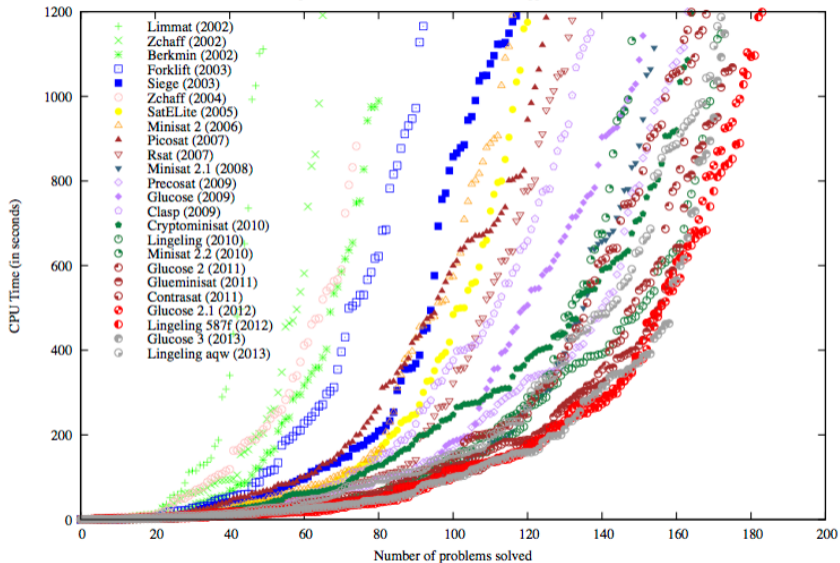
CDCL solvers have an expected polynomial $O(n^k)$ runtime on SAT instances whose primal (Gaifman) graph has treewidth k .

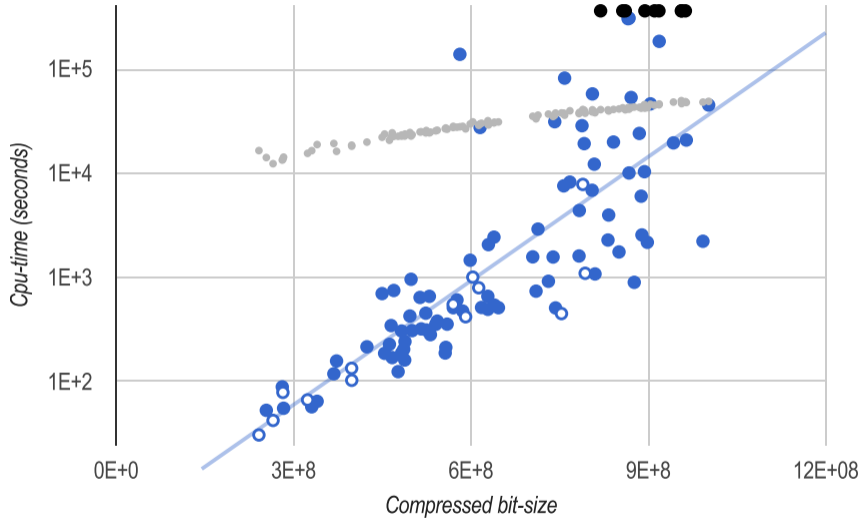
Without ever trying to compute a treewidth/decomposition (NP hard).

Go on a $n \times n$ goban is PSPACE-hard

- PSPACE-hard to decide if there is a winning strategy
- AlphaGo 0 does not solve 19×19 Go
- It plays better than humans (and that's amazing!)

Results of the SAT competition/race winners on the SAT 2009 application benchmarks, 20mn timeout





Additional ingredients (patented for some)

- (I) stops, restarts with a better understanding of the problem¹³
 - (I) forgets learnt information predicted as “useless” (Glue clauses²)
 - Lazy data structures²⁶
-
- Absolutely reliable combination of logic and intuition
 - but we don't really understand why it can be so efficient^{1,16}

Neural nets and safety critical settings

It doesn't seem too hard to fool a standard Convolutional Neural Net^a

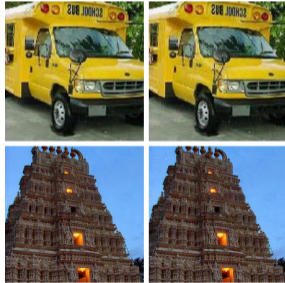
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Neural nets and safety critical settings

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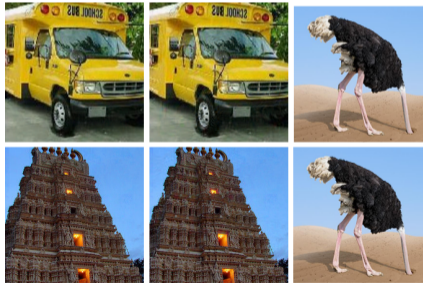
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Neural nets and safety critical settings

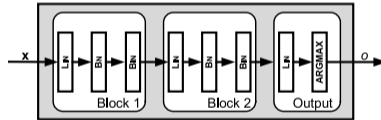
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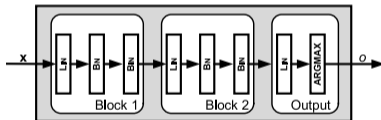
Binarized Deep NN: ± 1 activations/weights⁶

- Lin: affine transformation with learnt binary weights (float bias).
- Bn: (Batch normalization) rescaling with learnt floats.
- Bin: binarization using the *Sign* function.



Binarized Deep NN: ± 1 activations/weights⁶

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A learnt block can be described as a SAT^a formula

(SMT(LI)¹⁷ for ReLU)

^aNina Narodytska et al. "Verifying properties of binarized deep neural networks". In: *arXiv preprint arXiv:1709.06662* (2017).

Adversarial Robustness of a classifier

A positive test input cannot be slightly modified to change class

Adversarial Robustness of a classifier

A positive test input cannot be slightly modified to change class

Certified robustness by SAT

As a SAT formula: Neural Net + input + bounded perturbation + missclassification

Adversarial Robustness of a classifier

A positive test input cannot be slightly modified to change class

Certified robustness by SAT

As a SAT formula: Neural Net + input + bounded perturbation + missclassification

- MNIST dataset, 4 blocks BNN with 100 to 200 neurons per layer, L_∞ norm
- Millions of clauses: Glucose² certifies (non) robustness for most input in $< 5'$ CPU time

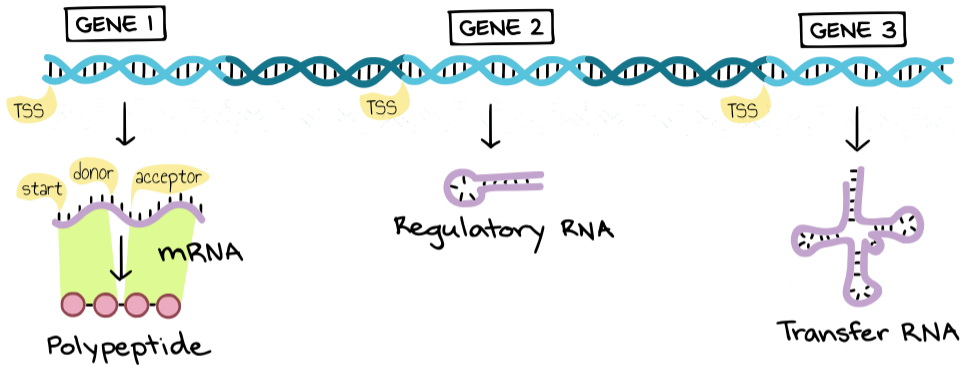


DNA

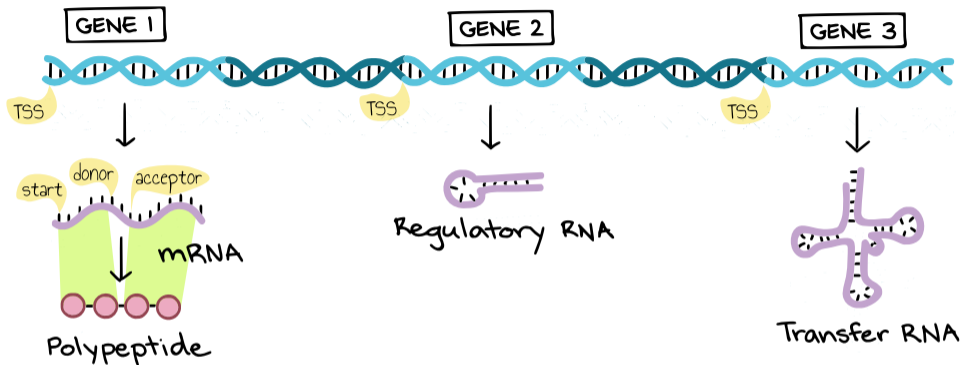
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Segmenting genomic DNA



Segmenting genomic DNA



Similarities

RNA-Seq
Rfam UniProt

Conservation



Sequence stats

Markov chains

Site predictors

SVM
Neural nets

Optimization + decomposable probability distribution

(Semi-CRF)²¹

- Derived from an actual human processor (S. Rumbauts, PhD)^a
- Discriminative learning (don't try to model evidence!)
- Optimizes an empirical loss function (performance on a testing set: quality is crucial)

^aS Foissac et al. "Genome Annotation in Plants and Fungi: EuGène as a Model Platform". In: *Current Bioinformatics* 3.2 (2008), pp. 87-97.



...

Prediction is in P

Main difficulty: collecting evidence, training and testing.